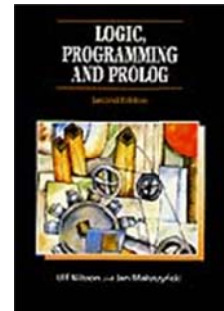


NOTATION IN NILSSON

Conventions seen in Chapter 1 and Appendix B of: Nilsson, Ulf and Małuszyński, *Logic, Programming and Prolog (2ed)*. <http://www.ida.liu.se/~ulfni/lpp/>



ROMAN ALPHABET

Description	Examples	typical use
lowercase italic letters	<i>a, b, c</i>	constants
	<i>f, g, h</i>	functors
	<i>p, q, r</i>	predicates
	<i>s, t</i>	terms
uppercase italic letters	<i>F, G</i>	formula
	<i>P</i>	set of premises (formulas)
	<i>S, T</i>	sets (of terms)
script capitals	<i>X, Y, Z</i>	variables
	\mathcal{A}	alphabet
	\mathcal{T}	set of terms
Fraktur/black-letter capitals	\mathcal{F}	set of (well-formed) functions
	\mathcal{D}	domain
	\mathfrak{I}	interpretation
double-struck capitals	\mathbb{N}	Set of natural numbers
	$\mathbb{Z}, \mathbb{Q}, \mathbb{R}$	Sets of integers, rationals, reals

Usage of italic and script letters can be arbitrary and inconsistent; they can be reassigned as convenient. Meanings of double-struck and fraktur letters are fairly standard and well-defined.

GREEK ALPHABET

Arial	Cambria	Symbol	CMMI	Letter name	Usage
φ	φ	φ	φ	phi	a valuation
θ	θ	θ	θ	theta	a substitution
σ	σ	σ	σ	sigma	a substitution; a valuation
γ	γ	γ	γ	gamma	a substitution
δ	δ	δ	δ	delta	a substitution
ε	ε	ε	ε	epsilon	the empty substitution

Different fonts can have confusingly different designs for Greek letters (at least, to my Roman-accustomed eyes). The typography in Nilsson uses “italic” forms of the Greek letters from Donald Knuth’s Computer Modern font series (CMMI). For example, φ (“phi”) tilts to the right and squishes a bit, becoming φ. (Greek typography doesn’t have true italics; instead, these seem to approximate handwritten forms.)

SYMBOLS

Glyph	Concept name	Example	Notes
Quantifiers			
\forall	universal quantifier	$\forall X$	“for all” or “for every”
\exists	existential quantifier	$\exists X$	“there exists some”
Logical connectives			
\wedge	conjunction	$x \wedge y$	“and”
\vee	disjunction	$x \vee y$	“or”
\neg	negation	$\neg x$	“not”
\supset	implication	$F \supset G$	“if... then...” (if F then G)
\leftarrow	implication	$G \leftarrow F$	“if... then...” (if F then G)
\leftrightarrow	equivalence		“if and only if” (“ iff ”)
\equiv	logical equivalence	$\neg \neg F \equiv F$	“is logically equivalent to”
\vdash	derivability	$P \vdash F$	“derives” or “ F is derivable from P ”
Sets			
\in	belonging	$x \in S$	“ x is in S ” or “ x is an element of S ”
\subseteq	subset (improper)		“subset of or equal to”
\cup	union	$S \cup T$	elements found in either
\cap	intersection	$S \cap T$	elements found in both
\times	Cartesian product	$\mathcal{D} \times \mathcal{D}$	
\emptyset	empty set		
$\{ \}$	set construction	$\{\langle \text{Adam} \rangle, \langle \text{Eve} \rangle\}$	
$\langle \rangle$	(angle brackets)	$p_{\mathfrak{S}} := \{\langle \text{Eve} \rangle\}$	<i>denotes individual vs. symbol?</i>
Miscellaneous			
	condition		“such that”
\circ	composition	$F \circ G$	
/	arity	p/n	
/	mapping (substitution)	$X/t \in \theta$	
$\frac{\text{premises}}{\text{conclusion}}$	predicate logic notation	$\frac{F \quad F \supset G}{G}$	“if F and F implies G , then G ”
$:=$	denotation, assignment	$zero := 0$	“zero denotes the number 0”
\mapsto	mapping (valuation)	$\varphi[X \mapsto t]$	
\models	truth	$\mathfrak{S} \models_{\varphi} Q$	“ Q is true with respect to \mathfrak{S} and φ ”
\models	logical consequence	$P \models F$	“ F is a logical consequence of P ”
$\not\models$	falsity	$\mathfrak{S} \not\models_{\varphi} Q$	“ Q is false with respect to \mathfrak{S} and φ ”