

Distributed Real-Time Combat Systems



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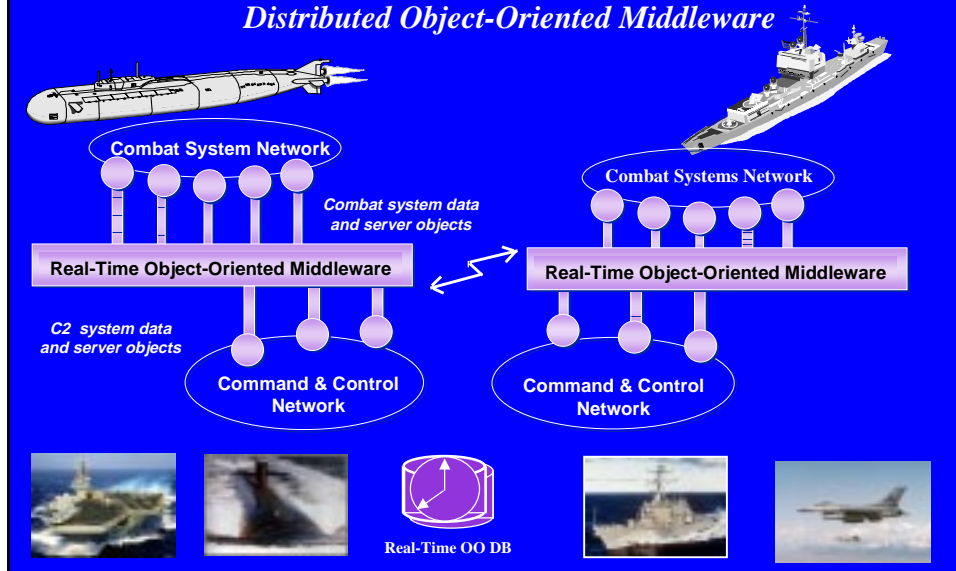
Operational Issues

- ◆ **Combat Systems and C2 Systems need to share data and functionality under *real-time constraints*.**
- ◆ **A need to maintain a consistent set of data at a specified level of QoS with the ability to enforce QoS tradeoffs in the middleware, databases, operating systems and networks.**
- ◆ **There is a need for heterogeneous data access, data sharing and data distribution across multiple platforms which may have different QoS requirements.**
- ◆ **Distributed systems need to meet specified timing constraints with the ability to negotiate for the QoS initially statically and then dynamically.**

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Operational Issues

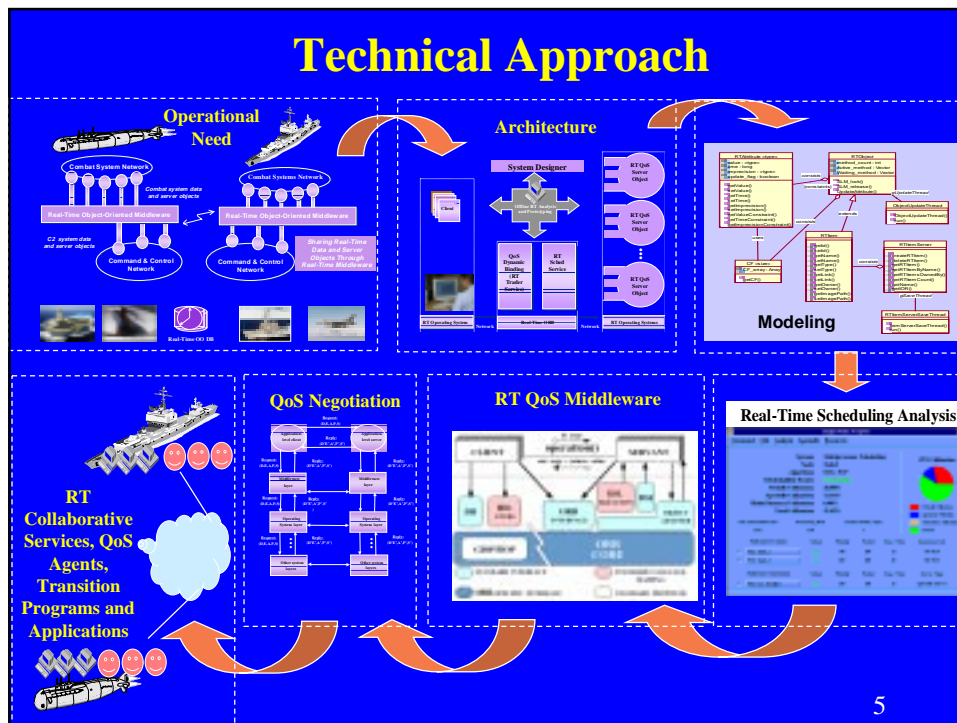
Real-Time QoS Data and Functionality Sharing Using Distributed Object-Oriented Middleware



Technical Approach

- ◆ Design real-time QoS model which enables the expression of time critical concepts and levels of QoS.
- ◆ Develop a multi-layered QoS negotiation schema which will provide a mathematical basis for synthesizing the parameters (real-time, accuracy, fault-tolerance & security).
- ◆ Develop a scheduling & analysis capability which provide metrics with respect to the synthesis of the parameters.
- ◆ Design a modular framework to support the system design and enable the insertion of custom solutions.
- ◆ Design real-time collaborative services which insure a consistent and current view of the data backed by guarantees and enforcement of timing constraints.

Technical Approach



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Accomplishments

- ◆ Implementation of the network-centric real-time QoS middleware algorithms and mechanisms. (static scheduling, dynamic scheduling, load shedding/reduction, dynamic binding, data replication)
- ◆ Implementation of QoS model in International Standard Unified Modeling Language (UML)
- ◆ Implementation of real-time QoS analysis tool with input from UML model and output to RT QoS middleware
- ◆ Implementation of QoS negotiations among real-time agents. (Accuracy vs Real-Time)
- ◆ Transitions:
 - Military programs (Coalition Forces, Virginia Class Sub, COE, Raytheon, Lockheed/Martin, Boeing, Mitre, TRW)
 - International standards (RT CORBA 1.0, RT CORBA 2.0, UML)
 - Commercial products (Analysis Tool, Scheduling Service, UML tools, WindRiver RTOS, Rational Software tools, Lineo Embedded Linux, OIS ORB)
 - Academic publications (IEEE TDPS, 2 Real-Time Systems Journal, conferences)

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Operational Payoff

- The ability for combat systems and C2 systems to share data and functionality under real-time QoS constraints.
- The ability to design and implement systems using a COTS middleware approach that many programs are adopting.
- New algorithms, mechanisms, and analysis techniques for distributed real-time QoS middleware.

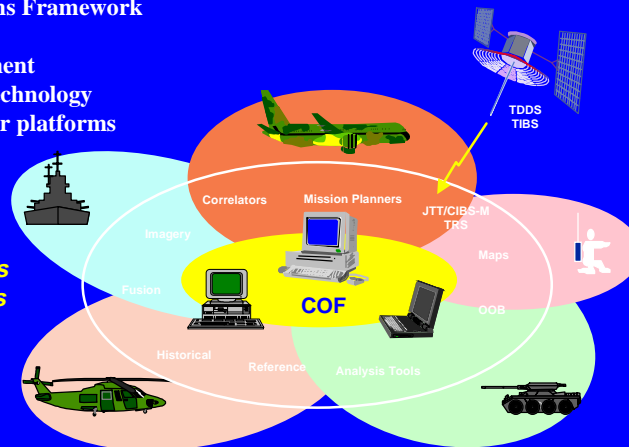
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Collaboration: Real-Time Support In Common Object Framework (COF)

Trudy Morgan, SPAWAR System Center SD

- Enhanced Data Integration - Object interfaces to legacy data types
- CORBA based Systems Framework
- COTS/GOTS Re-use
- Distributed Environment
- Extendable for new technology
- Transportable to other platforms

DRCS Real-time QoS middleware enforces real-time constraints in COF CORBA framework

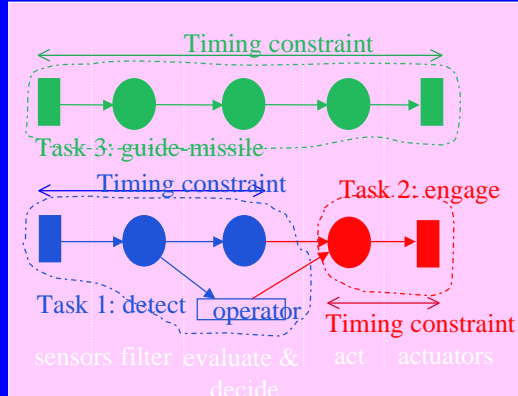


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Collaboration: Adaptive Resource Management in Asynchronous Real-Time Distributed Systems

Binoy Ravindran, Virginia Tech University

- ◆ Real-time computer systems for mission management
- ◆ Significant run-time uncertainties
 - Execution times, communication delays, event arrivals, etc.
- ◆ Require decentralization
 - Distributed application resources
 - Survivability
 - Meeting response time requirements



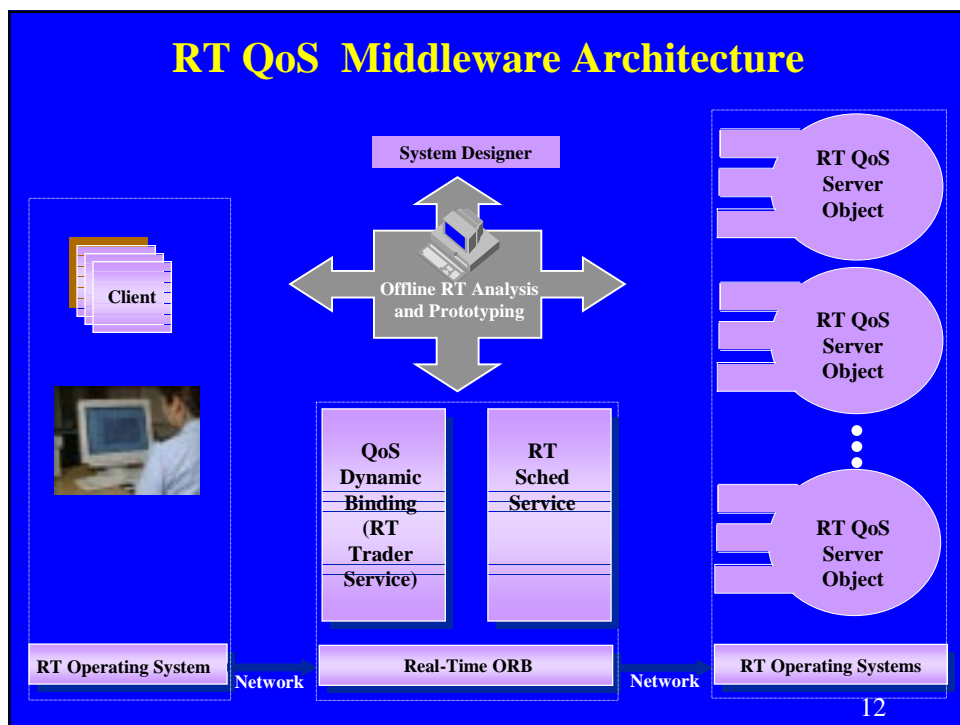
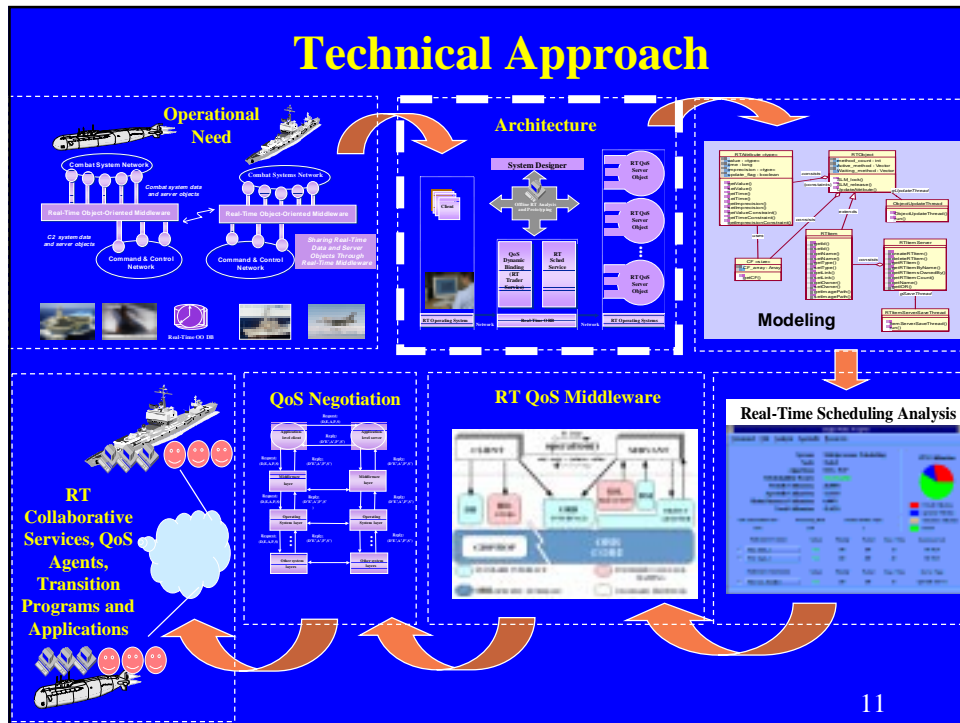
DRCS real-time QoS middleware algorithm and methodology sharing

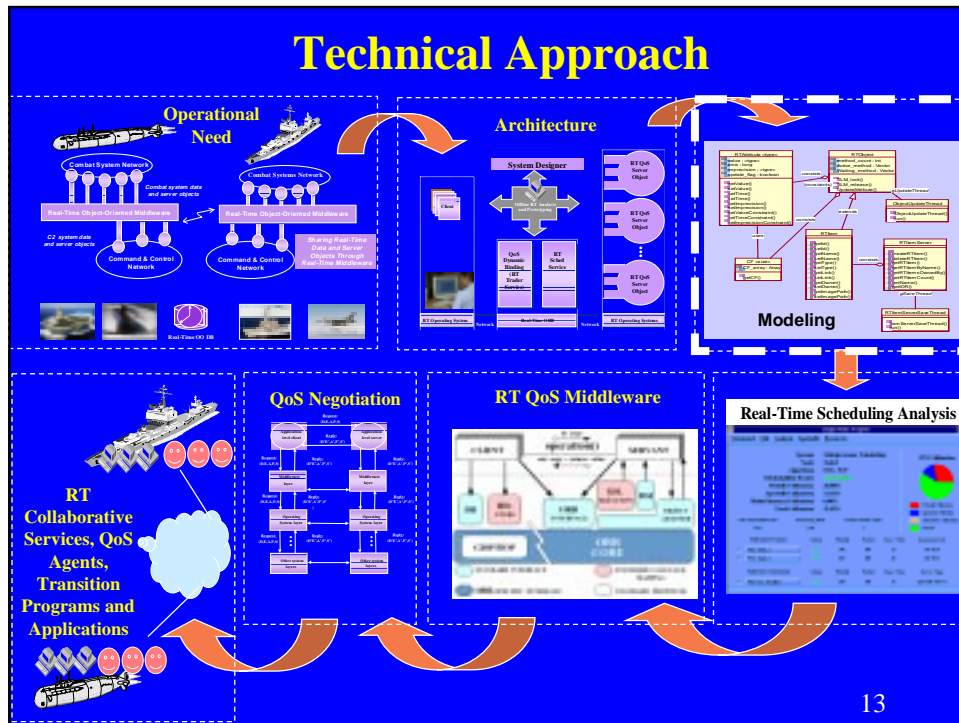
Technical Solution



Dr. Victor Fay-Wolfe

University of Rhode Island
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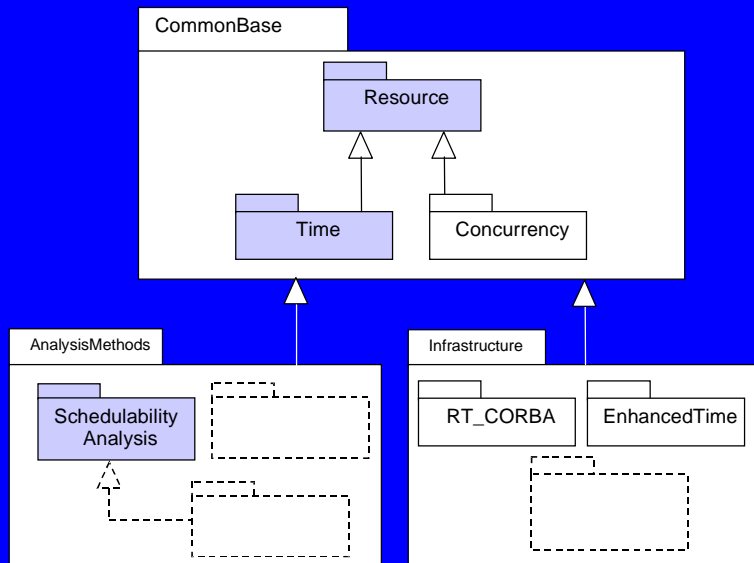
- ## RT QoS Modeling
- ◆ UML modeling of timing and QoS constraints
 - ◆ UML modeling of real-time QoS objects
 - ◆ Analysis modeling of RT QoS middleware
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QoS in Real-Time UML

- ◆ Responded to OMG Request for Proposal For RT Unified Modeling Language (UML)
- ◆ Not an extension to the UML metamodel, but a set of domain profiles for UML
- ◆ Goals
 - Enable the construction of models that could be used to make quantitative predictions regarding the characteristics of schedulability, performance and time
 - Facilitate communication of design intent between developers in a standard way
 - Enable interoperability between various analysis and design tools

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Profile Domain Packages



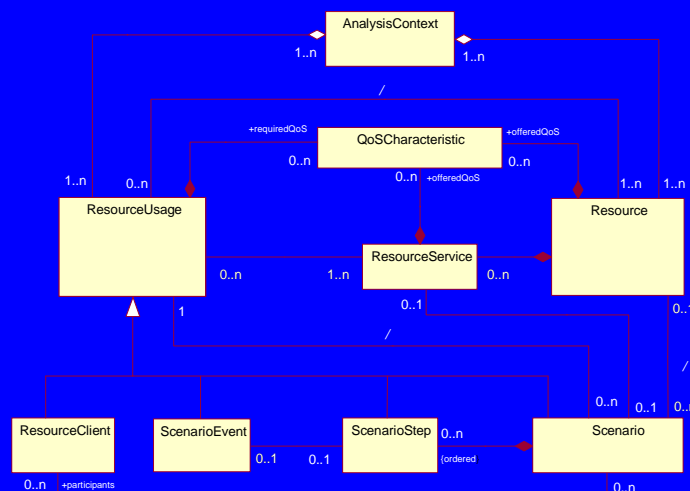
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Modeling Resources - The QoS Framework

- ◆ **Resource**: a model element that has some finite properties
 - reflects some finite physical quantity
 - may be logical (e.g., buffers) or physical
 - resources offer services (client-server model)
 - need to quantify the demand/supply of services
- ◆ **Quality of Service (QoS)**: a (usually quantitative) specification of:
 - the level of service required by a client from a resource or
 - the level of service offered by a resource to its clients

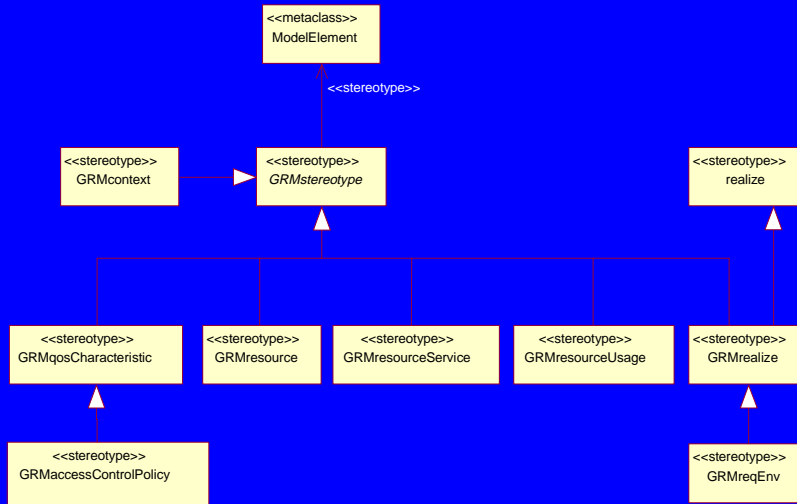
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The General Resource Model



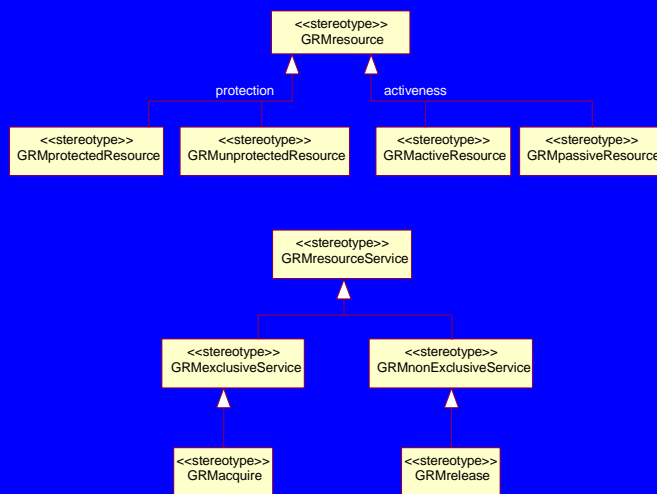
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Profile Stereotypes for Resource Model



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Profile stereotypes (cont.)



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Sample Schedulability Analysis (SA) Profile

SAEvent:

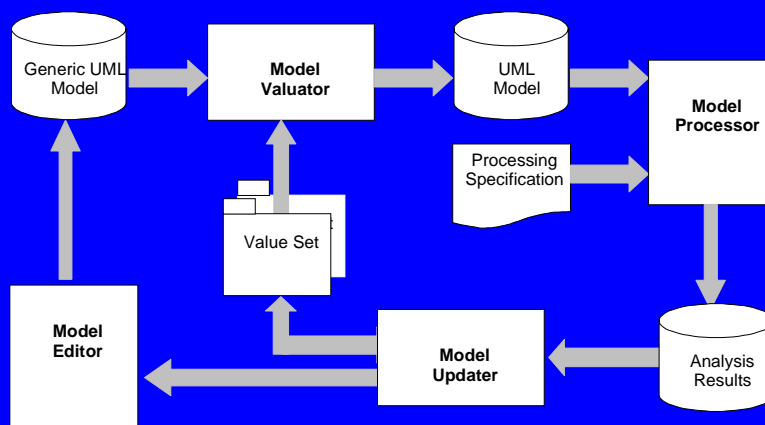
SARelativeDeadline = 400 ms

SAInstances = 1

SAOccurrencePattern = "exponential" 625

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Model Processing Framework



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Why Apply RMA to UML?

- ◆ UML addresses system structure and function
 - Multiple views of the system
 - Encourages top-down design
 - Ignores timing characteristics
- ◆ RMA addresses system timing characteristics
 - Uses the same system structure
 - Ignores the functional characteristics
- ◆ Two abstractions for the same system
 - Same structure
 - Different thinking

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The Usual Approach

- ◆ Focus on function
 - Meet the functional requirements
 - Timing is a requirement, but is too difficult to address
- ◆ Timing issues are addressed late in the process
 - Usually during system integration
 - The symptom ... sporadic failures
- ◆ Architectural changes are very expensive
 - Most timing problems require architectural changes
 - Usual consequence is an “over engineered” system

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A Better Approach

- ◆ Integrate timing considerations throughout the process
 - Start early in the design
 - Refine and update timing in concert with functional refinements
- ◆ Make timing specifications visible in the UML
 - Removes the dichotomy of functional and timing abstractions
 - Let tools construct the timing model
- ◆ Require timing validation during design, unit test, and integration

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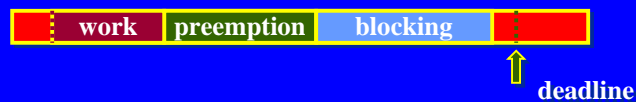
What Does RMA Reveal?

- ◆ RMA is static analysis
 - Not dynamic
 - Not simulation
- ◆ RMA establishes a bound for schedulability
 - A system is guaranteed to be schedulable within the bound
 - A system may run outside the bound, but is not guaranteed
- ◆ RMA shows how system resources are used
 - CPU and other “active” resources
 - Passive resources
 - Physical and logical resources

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When is a Task Schedulable?

- ◆ A task is schedulable if its worst case completion time is less than its deadline.
- ◆ Worst case completion time accounts for three classes of time
 - Work - execution of the task itself
 - Preemption - execution of higher priority tasks
 - Blocking - execution by lower or equal priority tasks when the task is ready for execution



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Analysis Results

- ◆ Schedulability
 - System, node, and task level
 - Utilization
 - Worst case completion times
- ◆ Analysis quality
 - Spare capacity
 - Blocking - total and by resource
 - Stability - behavior in overload
- ◆ “What if” analysis tools
 - The tool suggests changes
 - Rapid assessment for timing and architectural changes

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More Than Just RMA Analysis

- ◆ The analysis alone is very useful
 - Provides interesting information about a system
 - Enables “Timing Design”
- ◆ We can apply the analysis to another aspect of the problem
- ◆ We can answer the question, “How do I assign threads and priorities?”
 - Assign message priorities
 - Assign execution priorities
 - Assign capsules/activities to physical threads

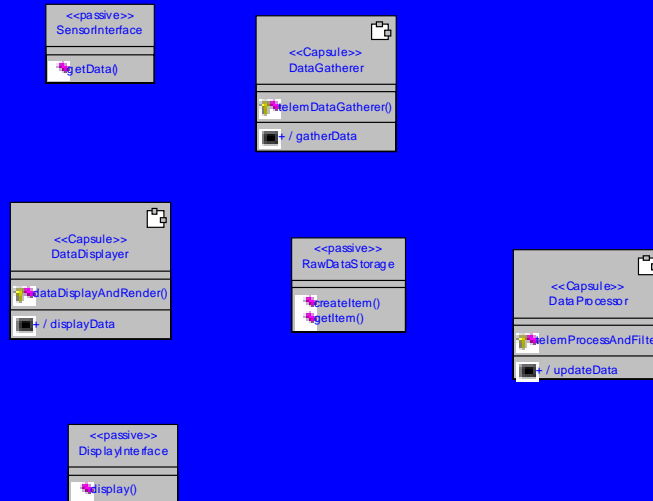
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Example System

- ◆ Telemetry System
 - Takes real-time data from a set of sensors
 - Filters and processes the data
 - Displays the filtered data to operator
- ◆ Display must be updated every 60 ms
- ◆ Telemetry data must be gathered every 100ms
- ◆ Filtered data update interval is 200ms

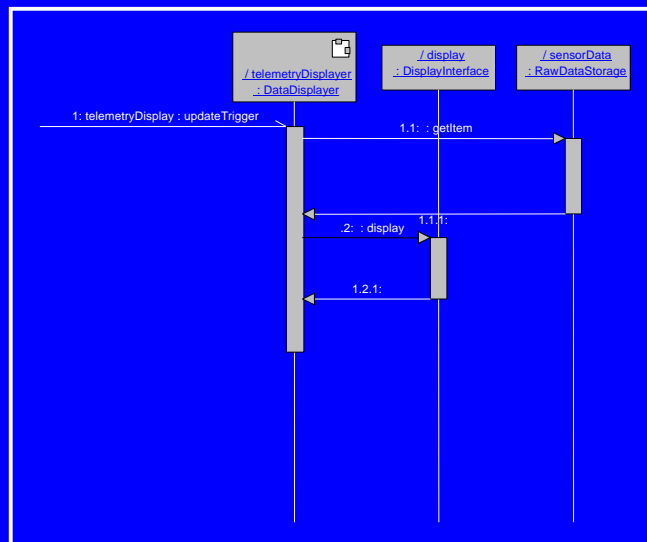
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RT UML Example: Class Diagram



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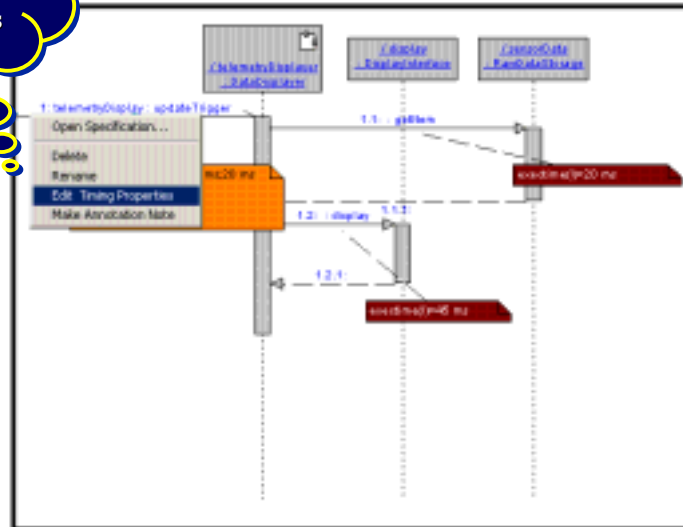
RT UML Example: Sequence Diagram



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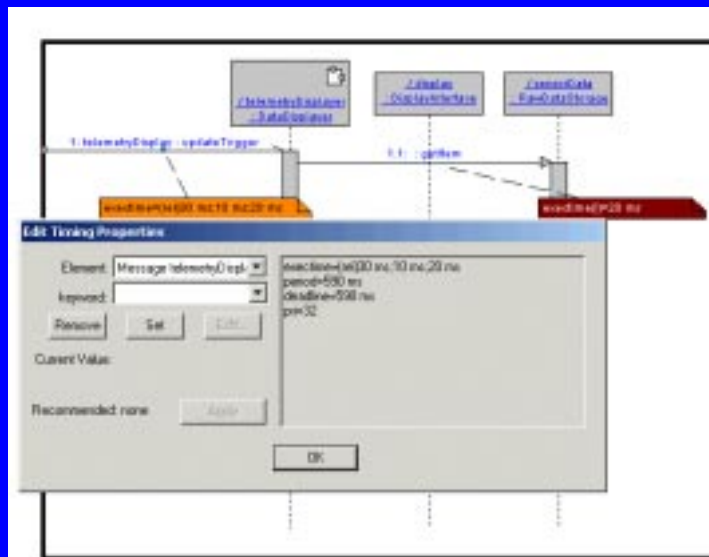
Sequence Diagram; Edit Timing Properties

Select any element; popup this menu



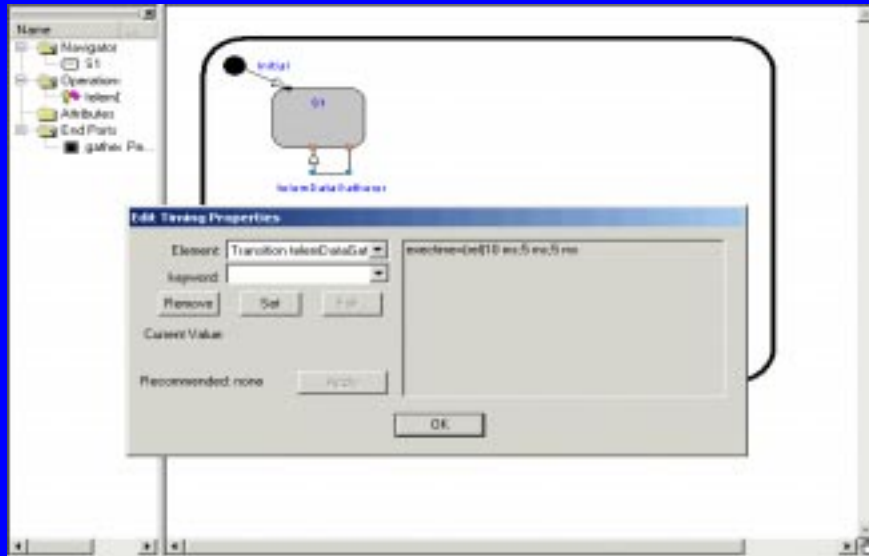
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Edit Timing Properties Dialog



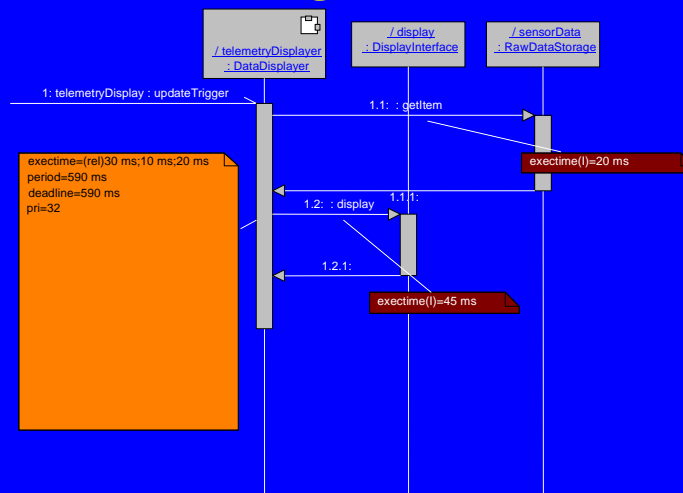
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Timing Properties on Transitions



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RT UML Example: Annotated Sequence Diagram



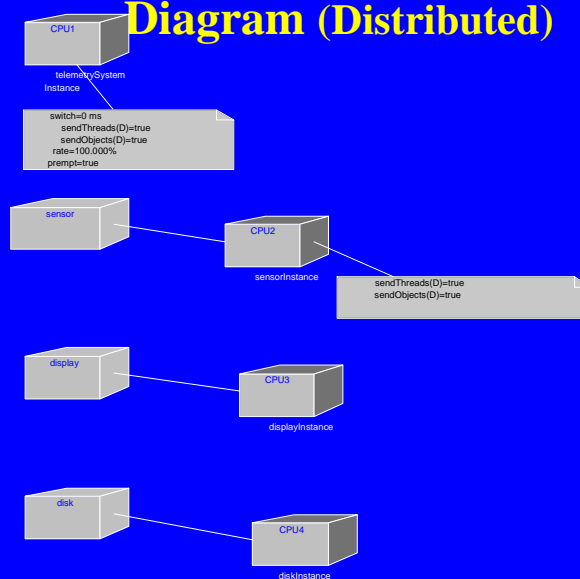
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RT UML Example: Distributed System

- ◆ You can design a distributed system with Rose RealTime
 - Use a deployment diagram to identify the processors
 - Assign package instances to specific processors
- ◆ DRCS tools can analyze a distributed system
 - Use end-to-end analysis or ROSA
- ◆ Network latency is expressed as another annotation on messages in the Sequence Diagram

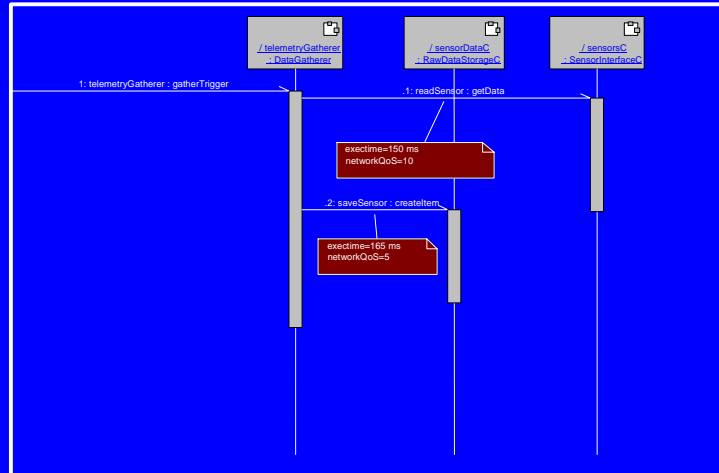
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RT UML Example: Deployment Diagram (Distributed)



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RT UML Example: Sequence Diagram (Distributed)



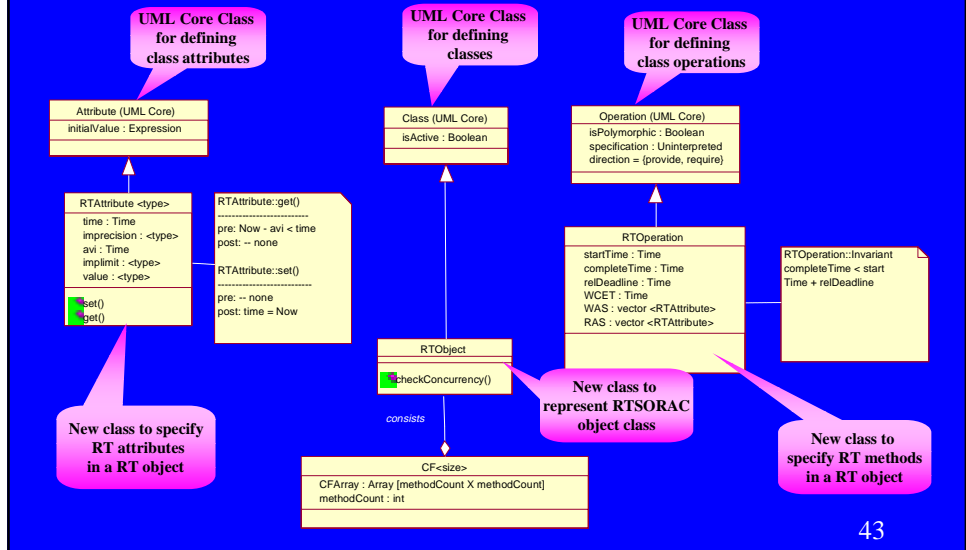
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Real-Time Objects

- ◆ **Definition:** A real-time object is an object in the environment that must be updated periodically to remain valid
- ◆ **Examples:**
 - radar or sonar readings
 - weather data

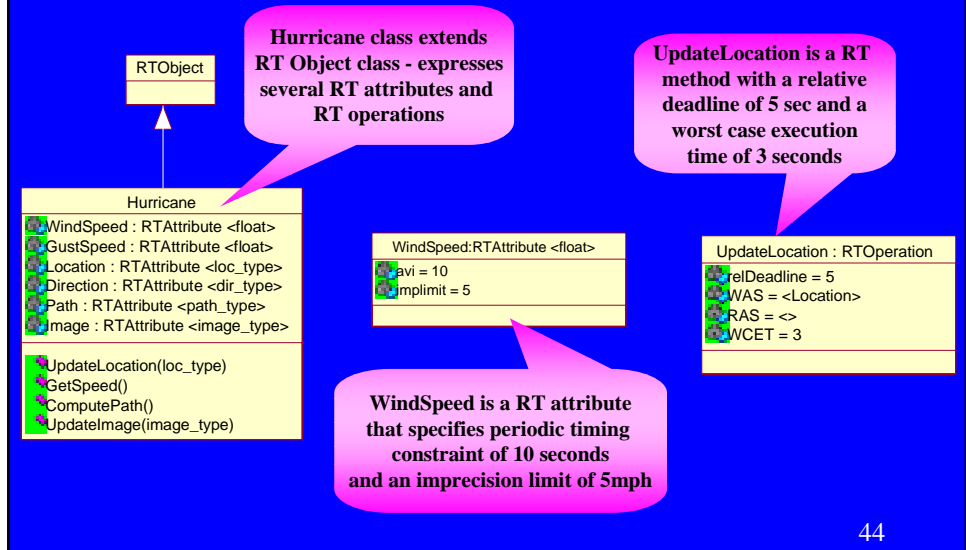
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UML Design for Real-Time Objects



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Example Hurricane Object



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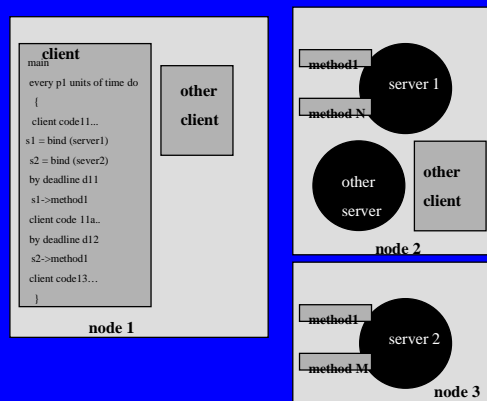
Modeling RT QoS Middleware

- ◆ A Real-Time QoS Middleware system is modeled from the client point of view.
- ◆ A client is partitioned into a sequence of dependent tasks.
- ◆ Middleware Services are modeled as resources.

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Modeling Real-Time QoS Middleware Application

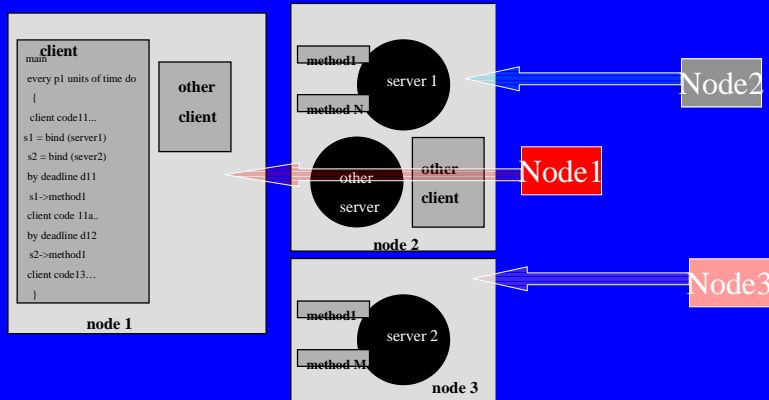
- ◆ Clients and Servers



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Modeling Real-Time QoS Middleware Application

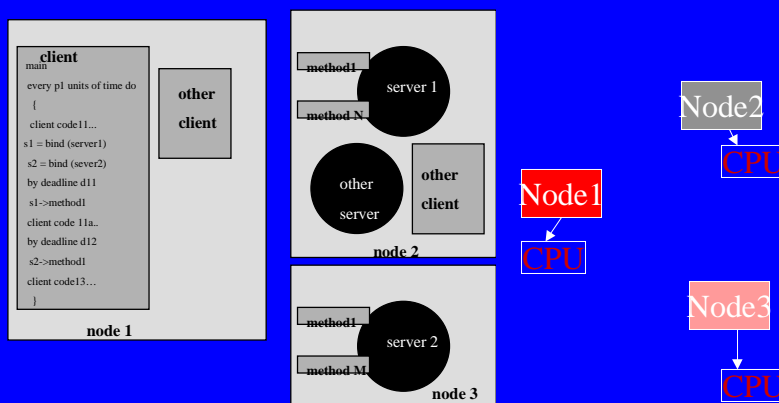
- ◆ All Servers as well as ORB and Services are represented by *resources*.



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Modeling Real-Time QoS Middleware Application

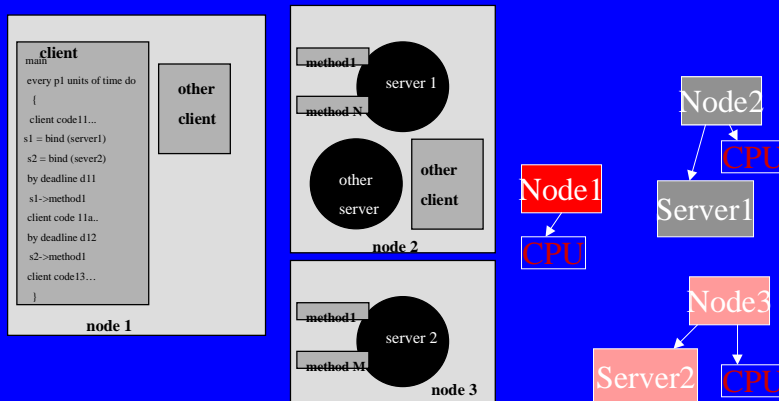
- ◆ Each node has a CPU as the processor.



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Modeling Real-Time QoS Middleware Application

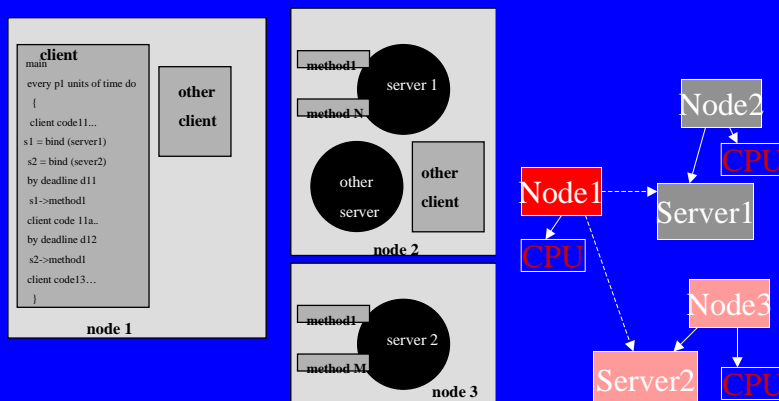
- ◆ The servers are represented as resources.



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Modeling Real-Time QoS Middleware Application

- ◆ The dotted arrows indicate access to the servers from the client.



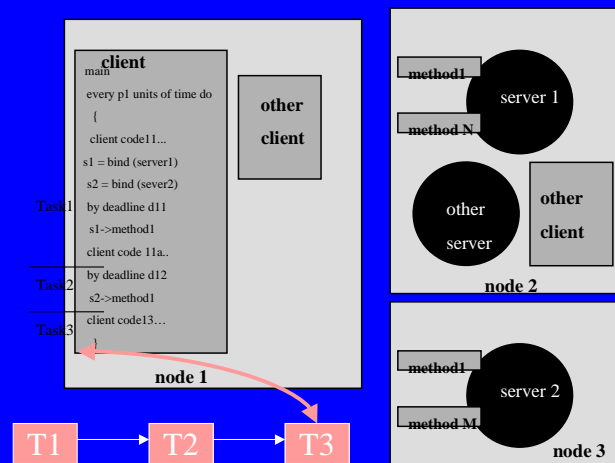
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Modeling Real-Time QoS Middleware Application

- ◆ CORBA Clients **can't** be mapped directly to *tasks* since Rate Monotonic Analysis does not support the analysis of the
 - *Intermediate Deadlines,*
 - *Network Delay.*
- ◆ Each Client with N intermediate deadlines will be modeled as N+1 dependent *tasks*.


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Modeling Real-Time QoS Middleware Application



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Modeling Real-Time QoS Middleware Application

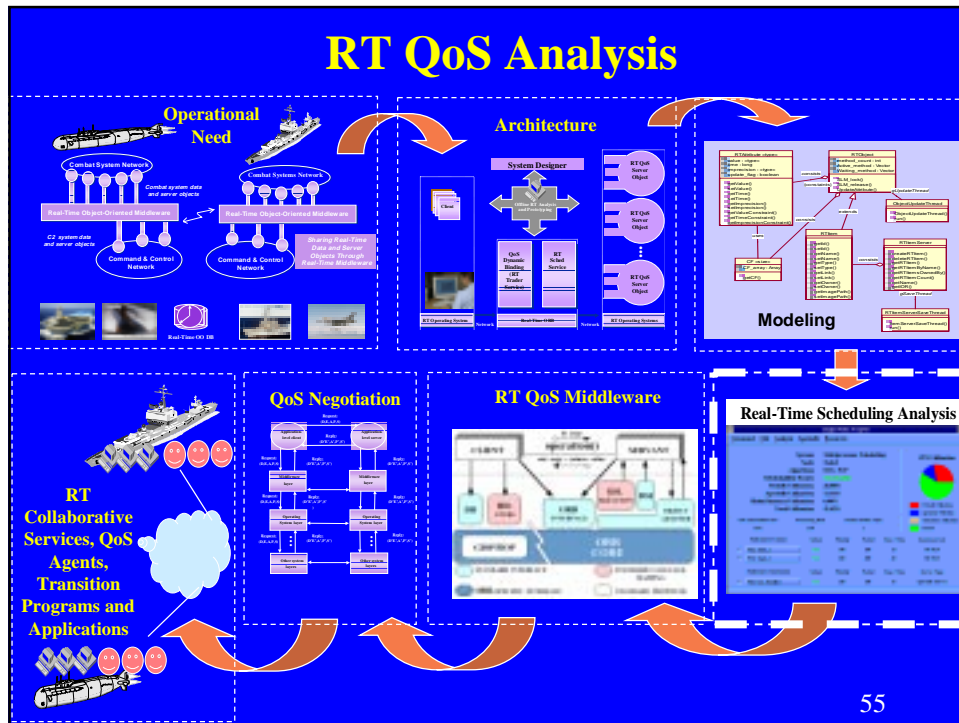
- ◆ Analysis: *END-TO-END* 
 - enables task dependencies analysis
- ◆ Priority assignment mechanism: *Deadline Monotonic (DM)*
 - Shorter the relative deadline - higher the priority.
- ◆ Resource access protocol: *DASPCP*
 - deadlock-free
 - limited blocking time
 - include network delay

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Modeling Real-Time QoS Middleware Application

- ◆ Analysis assumes that a system has 32K priority levels (consistent with RT CORBA standard).
- ◆ Typical operating systems do not permit that many priorities.
- ◆ DRCS maps the priority system into the limited priority systems on the network.
- ◆ The mapping minimizes the resulting priority inversions.

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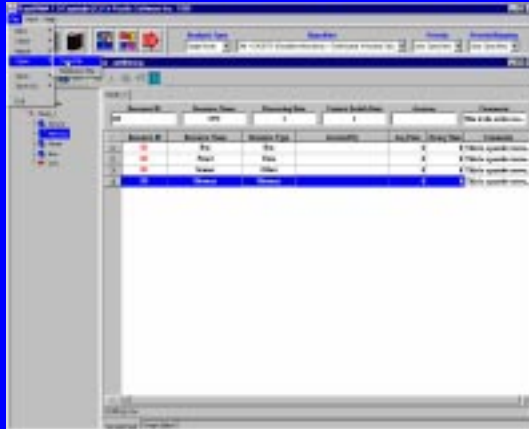


Modeling: Task Graph Editor

There are two components to our model. The task graph identifies the tasks of the system to be modeled.

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Modeling: Resource Graph Editor



The resource graph represents the resources in the system. This is a distributed system with global shared resources.

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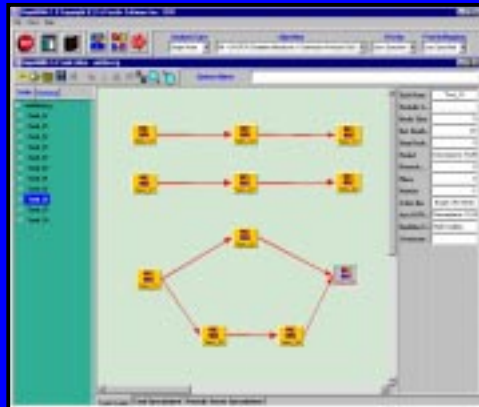
Modeling: Schedulability Analyzer



The Scheduler tests the system for **schedulability** (all tasks meet all their deadlines). There are a variety of scheduling analysis algorithms in the scheduler for various system architectures and scheduling methods.

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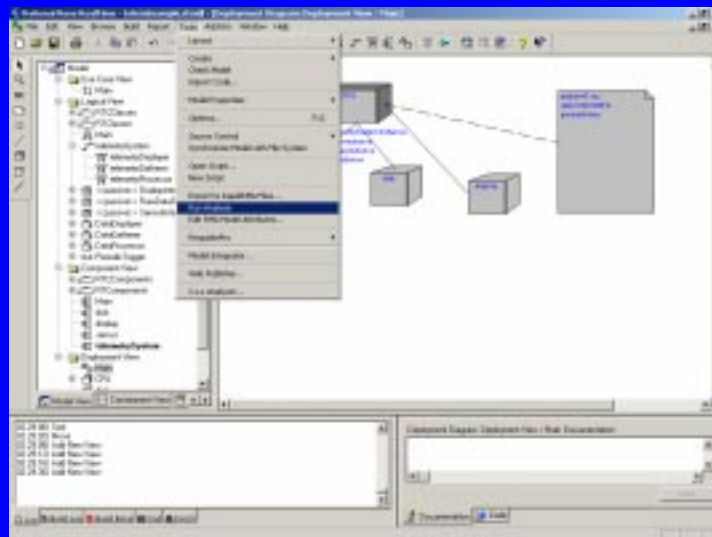
Modeling: End-to-End Analysis



End-to-End Scheduling analyzes the schedulability of a system with one or more paths of execution defined by a series of dependencies between tasks.

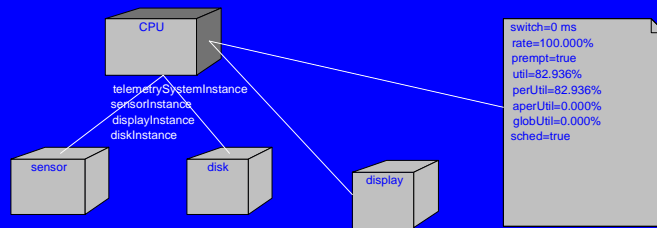
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RT UML Example: Run RMA Analysis



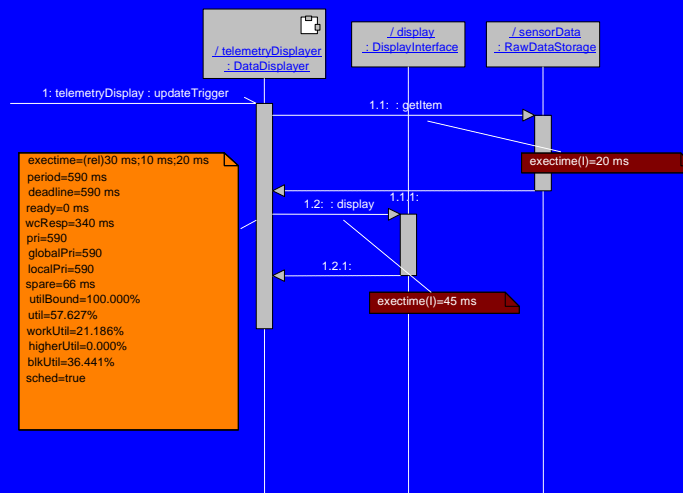
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RT UML Example: Analysis Results

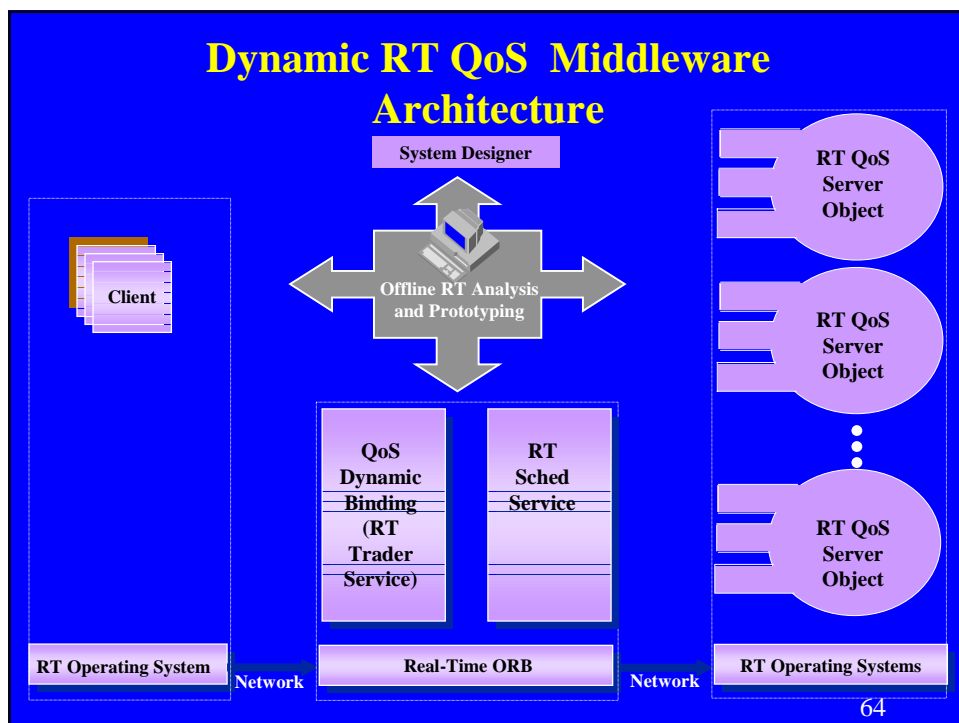
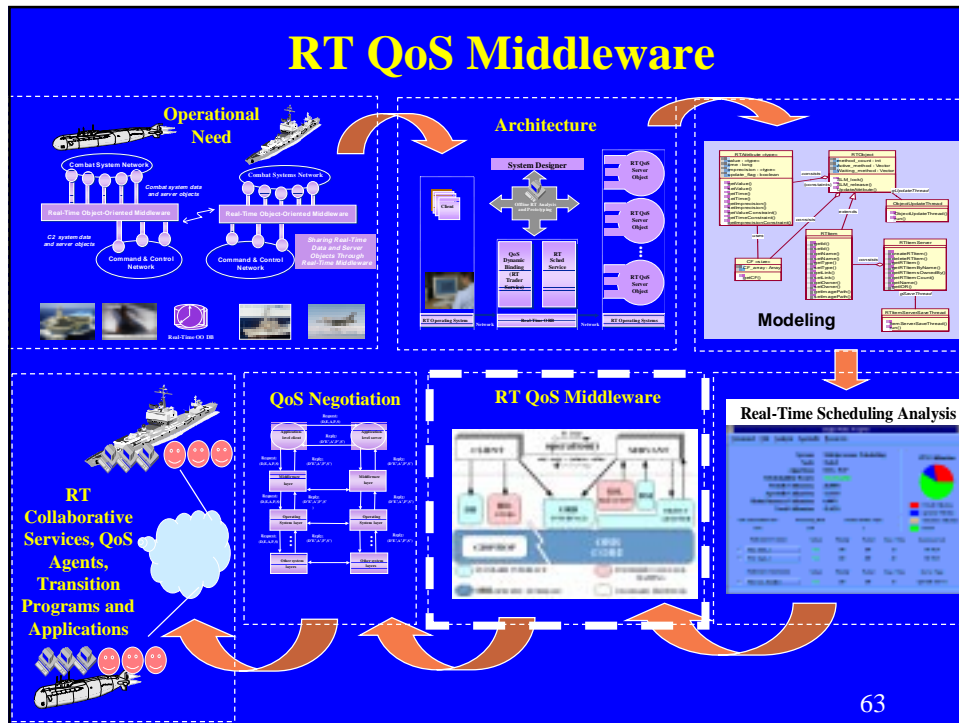


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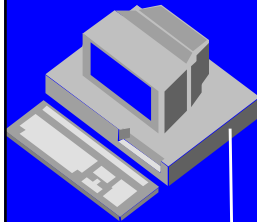
RT UML Example: Analysis Results



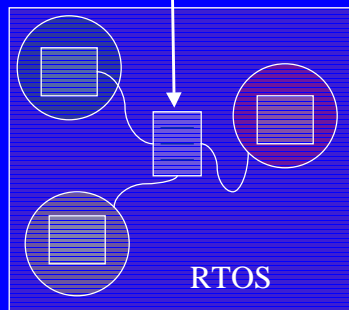
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Scheduling Service Architecture



*Generate output file with global priorities, local priorities, ceilings, etc.
This is the Scheduling Service configuration file.*

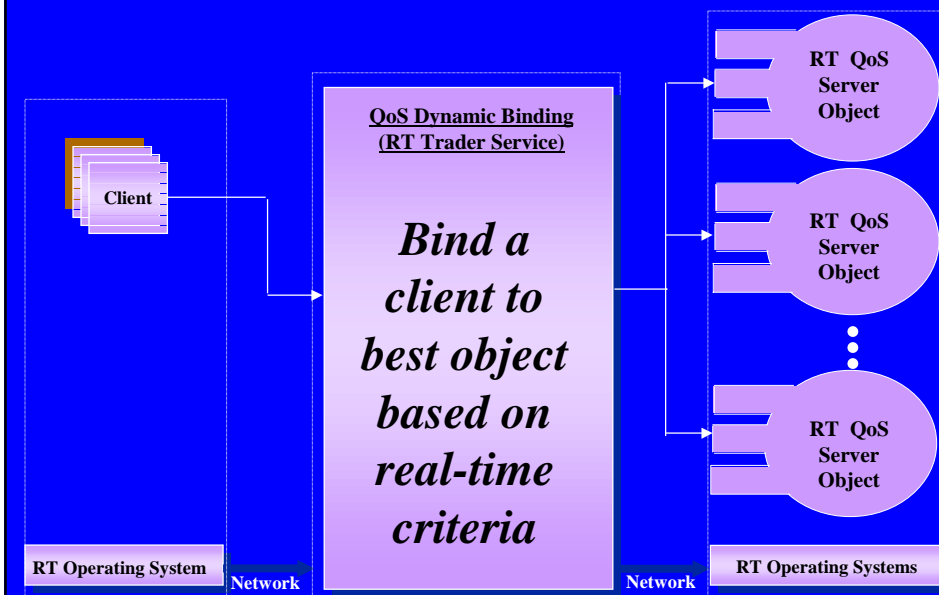


 Library code linked with every client and server

 Shared Memory Configuration file (global prios, local prios, ceilings)

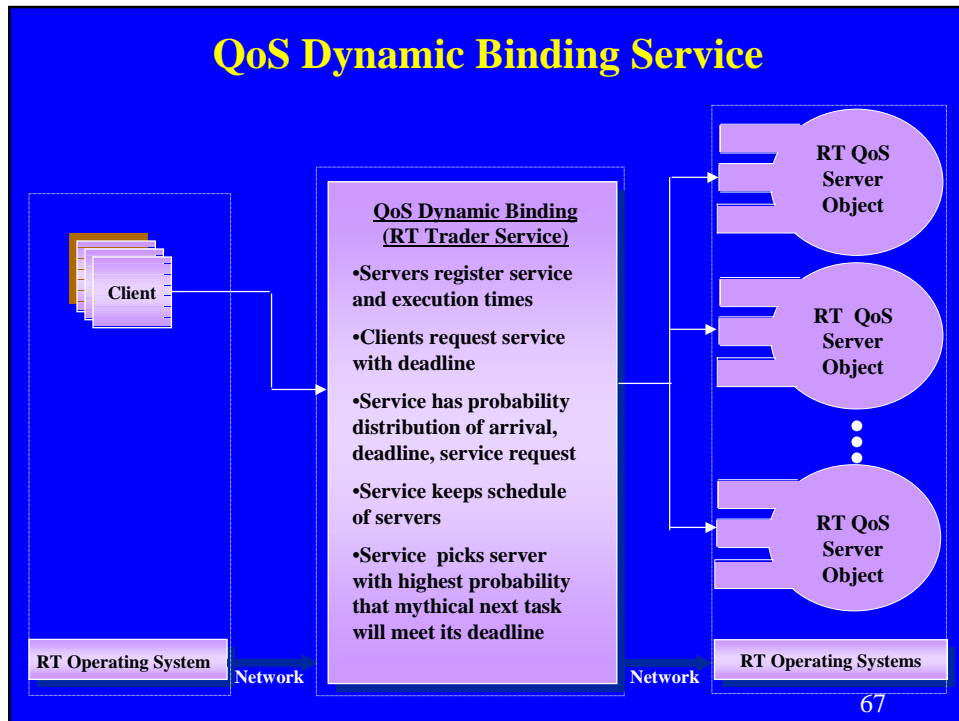
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QoS Dynamic Binding Service

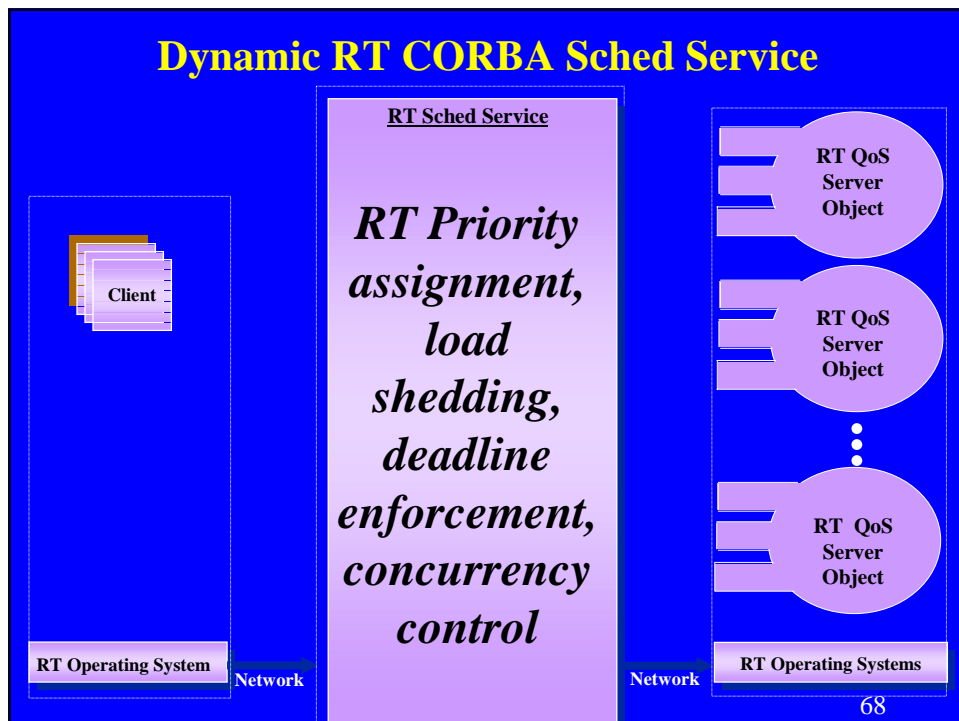


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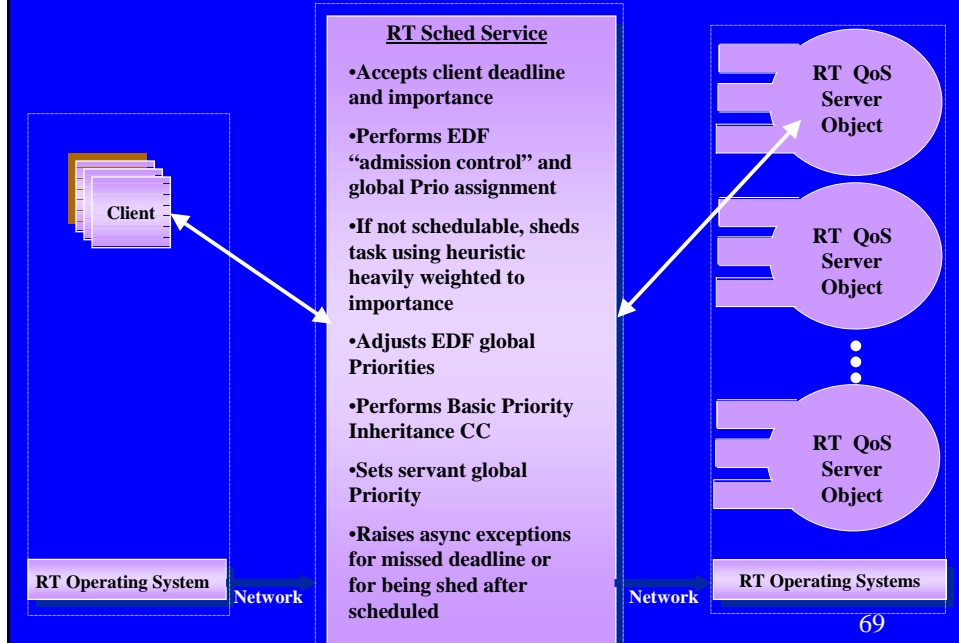
QoS Dynamic Binding Service



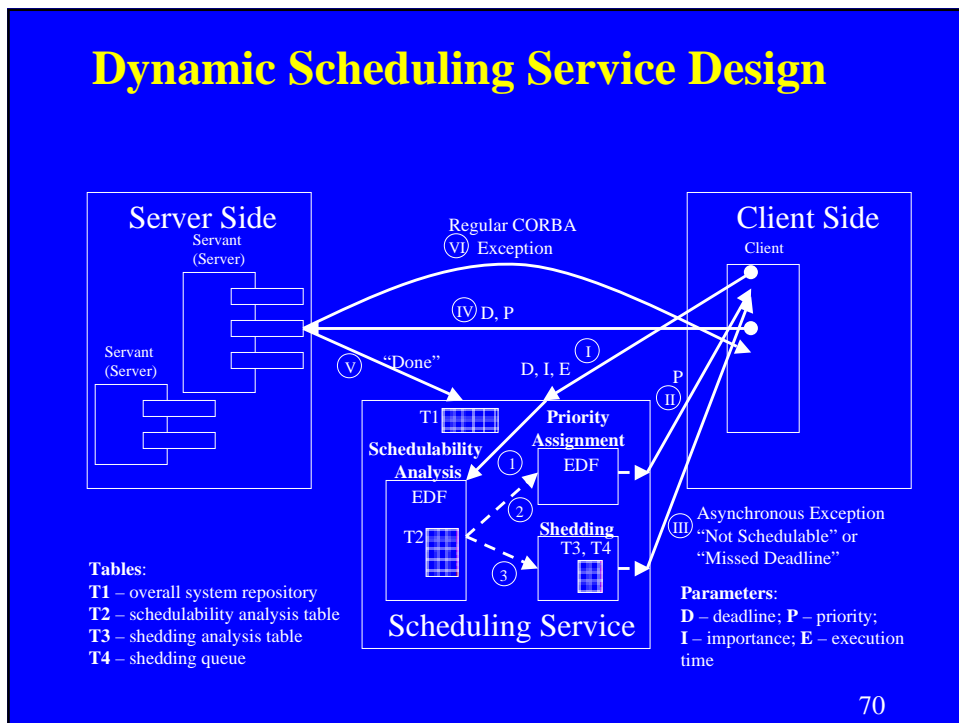
Dynamic RT CORBA Sched Service



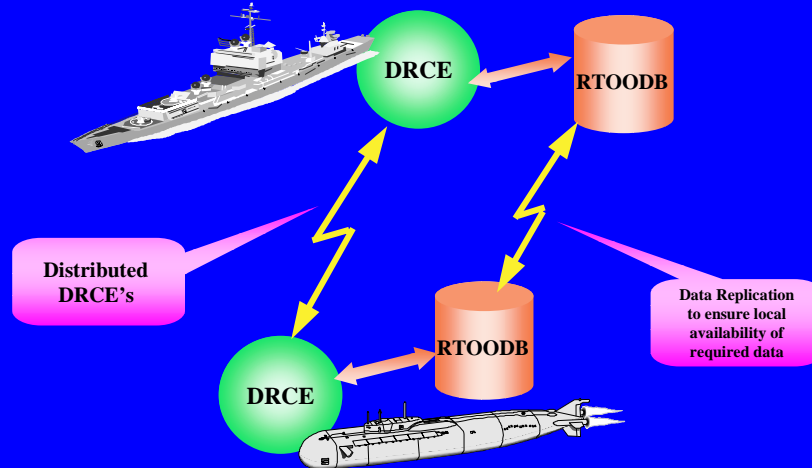
Dynamic RT QoS Middleware Scheduling Service



Dynamic Scheduling Service Design



RT QoS Middleware Databases



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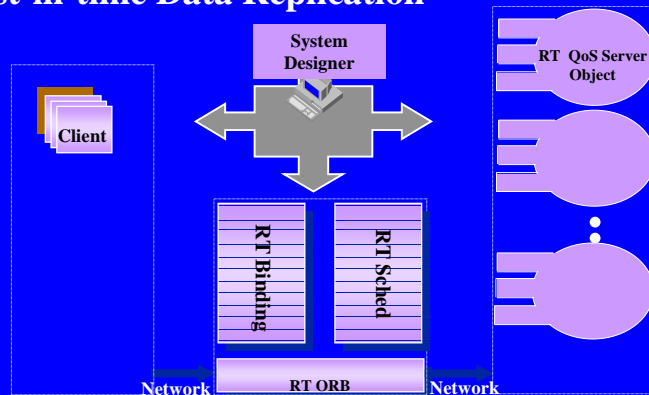
Real-Time QoS Object-Oriented Database Support

- ◆ Share data among collaborating users
 - possibly from remote sources
- ◆ Replicate data on remote sites
- ◆ Guarantee temporal validity of local copies of objects
- ◆ Just-In-Time Real-Time Replication Algorithms

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New Dynamic RT QoS Middleware Algorithms

- Load shedding (admission control) heuristics
- Dynamic priority mapping
- Affected Set Basic Priority Inheritance
- Just-in-time Data Replication



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Distributed Affected Set Priority Ceiling Algorithm

- ◆ *The conflict priority ceiling of a method m is the highest priority client that will ever lock a method that is not compatible with m ; where compatibility is defined by affected set semantics.*
- ◆ **Typical Priority Ceiling alg. steps used:**
 - grant lock if requesting priority > priority ceilings of all held locks.
 - use priority inheritance to reduce and bound blocking of high priority clients

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Distributed Affected Set Priority Ceiling Algorithm Properties

- ◆ *Consistency* - serializable object operations
- ◆ *Tight Priority Inversion Bound*
 - Proof similar to original Priority Ceiling results due to structure of protocol being the same, but granularity and conflict definition changed.
- ◆ *Deadlock prevention* - similar to previous Priority Ceiling results
- ◆ *Higher concurrency* - less blocking than in original Priority Ceiling algorithms
- ◆ *Efficient Implementation* - compatibility captured in Priority Ceiling check!

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Priority Mapping Problem Definition

- ◆ Real-Time CORBA 1.0 standard allows 32,000+ “CORBA priorities”
- ◆ RT OS have limited number of priorities
 - e.g. VXWorks, Lynx have 256 local priorities; Solaris 60
- ◆ RT middleware must map this large range of global priorities to RT OS priorities on heterogeneous nodes
- ◆ More than one global priority mapped to a local priority causes *priority inversion*
- ◆ Priority inversion must be accounted for as additional blocking time for task in analysis

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Priority Mapping Algorithm Solution

- ◆ Algorithm identifies how many overlapping priorities on each node
- ◆ Starts with lowest global priority and tries to “squeeze” it with next lowest.
- ◆ Performs schedulability check that includes new priority inversion blocking.
- ◆ If schedulable, those two global priorities are mapped to the same local priority. If not, then next highest global priority is tried for “squeeze”

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Priority Mapping Heuristics

- ◆ We have proven Priority Mapping algorithm to be optimal
- ◆ However, solution is NP-hard and takes excessive execution time
- ◆ We have developed several heuristics that are fast and near-optimal

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Dynamic Load Shedding Algorithm

- Let j be the index of “new” (coming) task in the Analysis Table (T2).
- Let n be the number of entries in the Analysis Table (T2).
- $\forall k = j..n$ compute slack time

$$t_{sl} = D_k - (t_c + \sum_{i=1..k} ER_i + B)$$

Blocking time B is essentially a place holder. It should be considered later.

The entries in the analysis table above the new one ($1..j-1$) are schedulable.

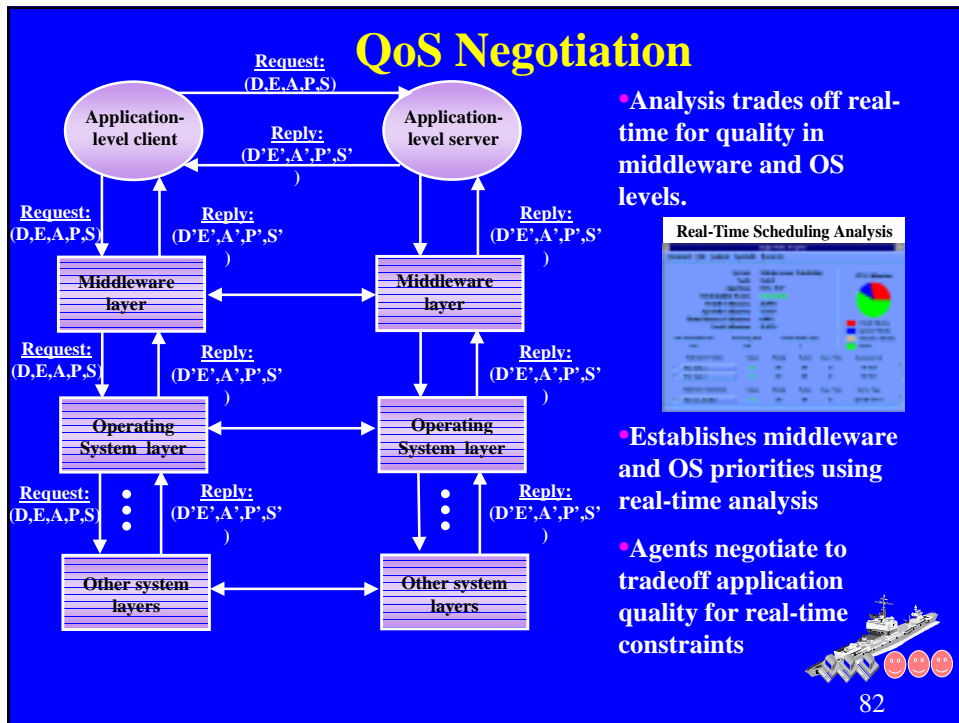
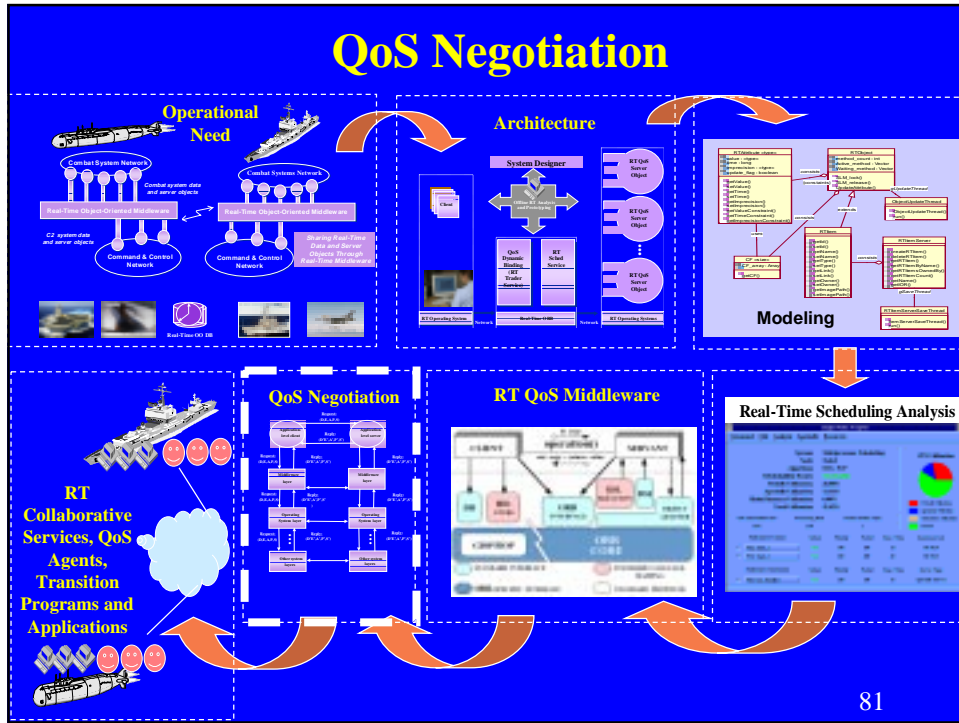
- If t_{sl} is negative then task k is unschedulable.

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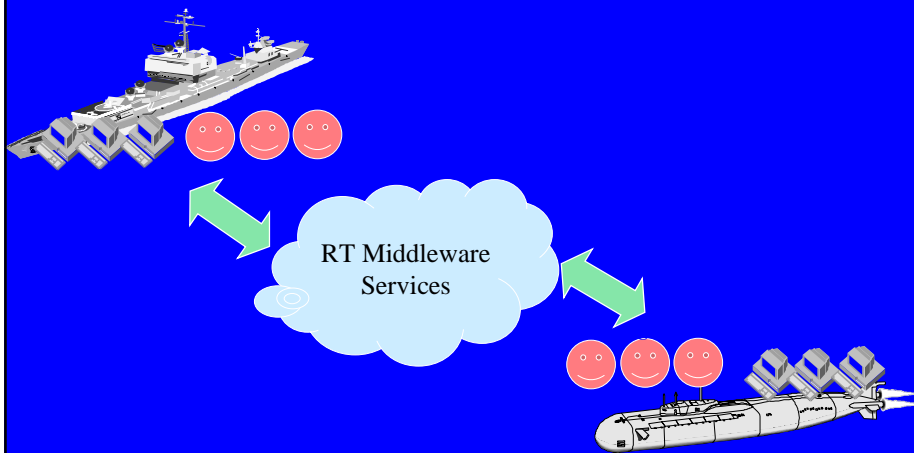
JIT-RT Data Replication Algorithms

- ◆ Static real-time environment
- ◆ Replication transactions
 - copy required data to local site
- ◆ Deadline computation
 - necessary and sufficient conditions for guaranteeing that all requests read temporally consistent data
- ◆ Replicates at two levels:
 - object level
 - method level - affected set semantics

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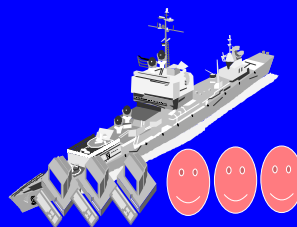
Distributed Real-time Collaborative Environment Architecture - RT Agents for QoS Negotiation



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Real-Time Agents

- ◆ **Definition:** A *real-time agent* is a flexible, autonomous software entity that must meet its design objectives within specified timing constraints.
 - model
 - architecture
 - communication
 - facilitation
 - scheduling



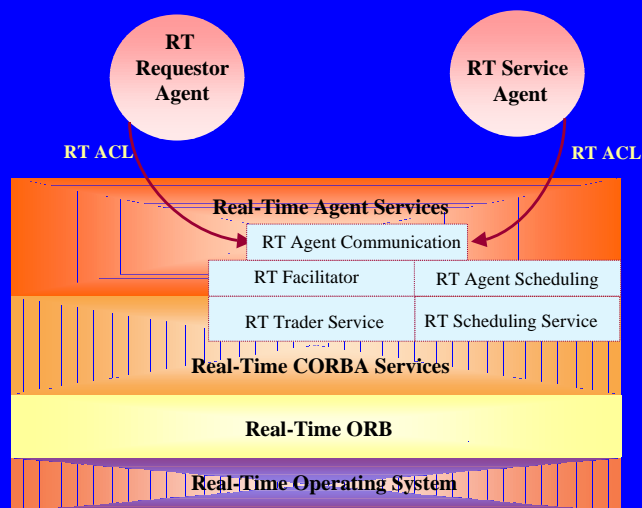
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Real-Time Agent Model

- ◆ **RT Agent**
 - solvables with multiple execution strategies
 - varying exec time and result quality
- ◆ **RT Agent Message**
 - deadline
 - importance
 - required quality

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RT Agent Architecture



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RT Agent Communication

- ◆ Extend agent communication language
- ◆ Express QoS within:
 - agent capabilities
 - agent requirements

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RT Agent Communication Language

```
(ask-one
 :sender      UserAgent
 :receiver    TrendWatcher
 :content     Watch(internet)
 :QoS_requirement (dl 15,imp 4,acc 75))
```

Ask a fellow agent
to watch a stock
trend within 15 seconds

```
(advertise
 :sender      BuyerSeller
 :receiver    Facilitator
 :content     BuyStock(A)
 :QoS_capabilities(
              (ex 5, acc 85)
              (ex 2, acc 65)))
```

Advertise QoS
capabilities to facilitator
including worst case
execution times and
accuracy levels

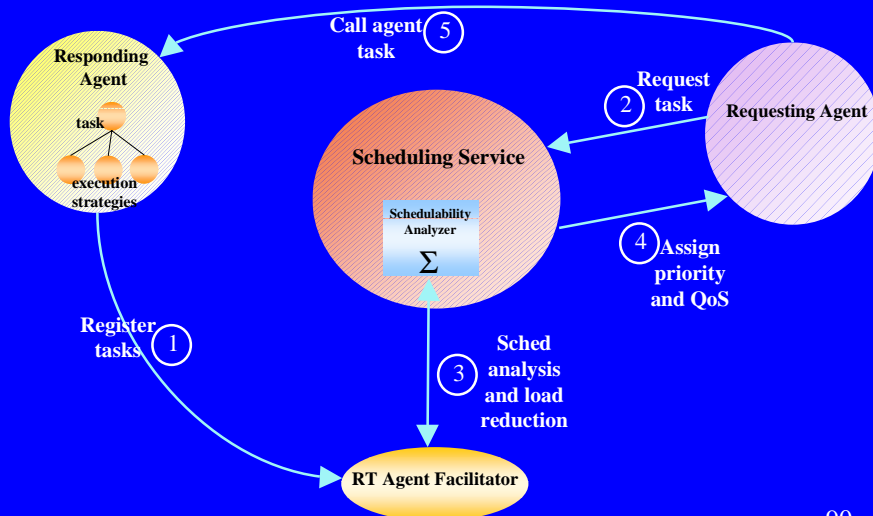
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RT Agent Scheduling

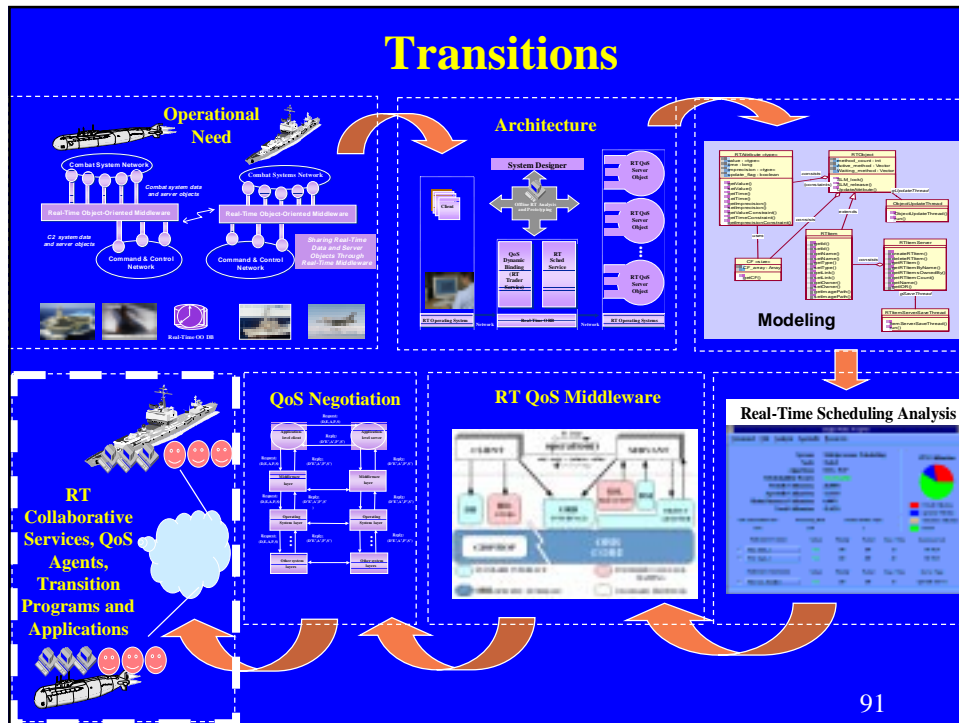
- ◆ Extend RT CORBA Load Shedding scheduling algorithm
- ◆ *Load Reduction*
 - if system of tasks cannot be scheduled
 - reduce in quality of one or more tasks to gain more execution time for schedule

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RT Agent Scheduling and Facilitation



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- # Accomplishments
- ◆ Implementation of the network-centric real-time QoS middleware algorithms and mechanisms. (static scheduling, dynamic scheduling, load shedding/reduction, dynamic binding, data replication)
 - ◆ Implementation of QoS model in International Standard Unified Modeling Language (UML)
 - ◆ Implementation of real-time QoS analysis tool with input from UML model and output to RT QoS middleware
 - ◆ Implementation of QoS negotiations among real-time agents. (Accuracy vs Real-Time)
 - ◆ Transitions:
 - Military programs (Coalition Forces, Virginia Class Sub, COF, Raytheon, Lockheed/Martin, Boeing, Mitre, TRW)
 - International standards (RT CORBA 1.0, RT CORBA 2.0, UML)
 - Commercial products (Analysis Tool, Scheduling Service, UML tools, WindRiver RTOS, Rational Software tools, Lineo Embedded Linux, OIS ORB)
 - Academic publications (IEEE TDPS, 2 Real-Time Systems Journal, conferences)

Operational Payoff

- The ability for combat systems and C2 systems to share data and functionality under real-time QoS constraints.
- The ability to design and implement systems using a COTS middleware approach that many programs are adopting.
- New algorithms, mechanisms, and analysis techniques for distributed real-time QoS middleware.

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Web Sites



atticus.spawar.navy.mil/dhda



homepage.cs.uri.edu/research/rtisorac/



www.tripac.com

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