Equivalence Relations

Section 9.5

Equivalence Relations

Definition: A relation on a set *A* is called an *equivalence relation* if it is reflexive, symmetric, and transitive.

Equivalence Relations

Definition: Two elements a, and b that are related by an equivalence relation are called *equivalent*. The notation $a \sim b$ is often used to denote that a and b are equivalent elements with respect to a particular equivalence relation.

Strings

Example: Suppose that R is the relation on the set of strings of English letters such that aRb if and only if l(a) = l(b), where l(x) is the length of the string x. Is R an equivalence relation?

Proof: Show that all of the properties of an equivalence relation hold.

- Reflexivity: Because l(a) = l(a), it follows that aRa for all strings a.
- Symmetry: Suppose that aRb. Since l(a) = l(b), l(b) = l(a) also holds and therefore bRa.
- Transitivity: Suppose that aRb and bRc. Since l(a) = l(b), and l(b) = l(c), l(a) = l(c) also holds and therefore aRc. (QED)

Relation on Real Numbers

Example: Let R be the relation on the set of real numbers such that aRb if and only if a – b is an integer. Is R an equivalence relation?

Proof: Because a - a = o is an integer for all real numbers a, aRa for all real numbers a. Hence, R is reflexive. Now suppose that aRb. Then a - b is an integer, so b - a is also an integer. Hence, bRa. It follows that R is symmetric. If aRb and bRc, then a - b and b - c are integers. Therefore, a - c = (a - b) + (b - c) is also an integer. Hence, aRc. Thus, R is transitive. Consequently, R is an equivalence relation. (QED)

Bit String Equivalences

- **Example**: Let n be a positive integer and S a set of strings. Suppose that Rn is the relation on S such that sRn t if and only if s = t, or both s and t have at least n characters and the first n characters of s and t are the same. That is, a string of fewer than n characters is related only to itself; a string s with at least n characters is related to a string t if and only if t has at least n characters and t begins with the n characters at the start of s. For example, let n = 3 and let S be the set of all bit strings. Then sR_3 t either when s = t or both s and t are bit strings of length 3 or more that begin with the same three bits. For instance, or R3 o1 and 00111 R3 00101, but 01 R 3 010 is not and 01011 R 3 01110 is not.
- Show that for every set S of strings and every positive integer n, R_n is an equivalence relation on S.

Bit String Equivalences

- **Solution**: We show that the relation Rn is reflexive, symmetric, and transitive.
 - *Reflexive*: The relation Rn is reflexive because s = s, so that sRn s whenever s is a string in S.
 - Symmetric: If sRn t, then either s = t or s and t are both at least n characters long that begin with the same n characters. This means that tRn s. We conclude that Rn is symmetric.
 - Transitive: Now suppose that sRn t and tRn u. Then either s = t or s and t are at least n characters long and s and t begin with the same n characters, and either t = u or t and u are at least n characters long and t and u begin with the same n characters. From this, we can deduce that either s = u or both s and u are n characters long and s and u begin with the same n characters, i.e. s Rn u. Consequently, Rn is transitive.
- It follows that Rn is an equivalence relation. (QED)

Divides

Example: Show that the "divides" relation on the set of positive integers is not an equivalence relation.

Solution: The properties of reflexivity, and transitivity do hold, but there relation is not symmetric. Hence, "divides" is not an equivalence relation.

- Reflexivity: a divides a for all a.
- *Not Symmetric*: For example, 2 divides 4, but 4 divides 2 does not hold. Hence, the relation is not symmetric.
- Transitivity: Suppose that a divides b and b divides c. Then there are positive integers k and l such that b = ak and c = bl. Hence, c = a(kl), so a divides c. Therefore, the relation is transitive.

Equivalence Classes

Definition: Let R be an equivalence relation on a set A. The set of all elements that are related to an element a of A is called the *equivalence class* of a. The equivalence class of a with respect to R is denoted by $[a]_R$.

When only one relation is under consideration, we can write [a], without the subscript R, for this equivalence class.

Note that $[a]_R = \{s \mid (a,s) \in R\}.$

• If $b \in [a]_R$, then b is called a representative of this equivalence class. Any element of a class can be used as a representative of the class.

Bit String Equivalence Class

- **Example**: What is the equivalence class of the string one with respect to the equivalence relation R₃ from the Bit String Equivalence example on the set of all bit strings? (Recall that sR₃ t if and only if s and t are bit strings with s = t or s and t are strings of at least three bits that start with the same three bits.)

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[011]_{R_3} = \{011, 0110, 0111, 01100, 01101, 01110, 01111, \ldots\}
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Equivalence Classes

Theorem 1: let *R* be an equivalence relation on a set *A*. These statements for elements *a* and *b* of *A* are equivalent:

- (i) aRb
- (ii) [a] = [b]
- (iii) $[a] \cap [b] \neq \emptyset$

Equivalence Classes

Proof:

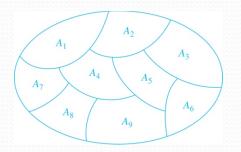
- (a) We show that (*i*) implies (*ii*). Assume that aRb. Now suppose that $c \in [a]$. Then aRc. Because aRb and R is symmetric, bRa. Because R is transitive and bRa and aRc, it follows that bRc. Hence, $c \in [b]$. Therefore, $[a] \subseteq [b]$. A similar argument shows that $[b] \subseteq [a]$. Since $[a] \subseteq [b]$ and $[b] \subseteq [a]$, we have shown that [a] = [b].
- (b) We show that (ii) implies (iii). Assume that [a] = [b]. It follows that $[a] \cap [b] \neq \emptyset$ since [a] is nonempty; $a \in [a]$ because R is reflexive.
- (c) We show that (iii) implies (i). Suppose that $[a] \cap [b] \neq \emptyset$. Then there is an element c with $c \in [a]$ and $c \in [b]$. In other words, aRc and bRc. By the symmetric property, cRb. Then by transitivity, because a c and cRb, we have aRb.

Because (i) implies (ii), (ii) implies (iii), and (iii) implies (i), the three statements, (i), (ii), and (iii), are equivalent.

Partition of a Set

Definition: A partition of a set S is a collection of disjoint nonempty subsets of S that have S as their union. In other words, the collection of subsets A_i , where $i \in I$ (where I is an index set), forms a partition of S if and only if

- $A_i \neq \emptyset$ for $i \in I$,
- $A_i \cap A_j = \emptyset$ when $i \neq j$,
- and $\bigcup_{i \in I} A_i = S$.



A Partition of a Set

An Equivalence Relation Partitions a Set

• Let R be an equivalence relation on a set A. The union of all the equivalence classes of R is all of A, since an element a of A is in its own equivalence class $[a]_R$. In other words, $[a]_R = A$.

• From Theorem 1, it follows that these equivalence classes are either equal ([a] = [b] with $[a] \cap [b] \neq \emptyset$) or disjoint ($[a]_R \neq [b]_R$ with $[a]_R \cap [b]_R = \emptyset$).

 $a \in A$

• Therefore, the equivalence classes form a partition of *A*, because they split *A* into disjoint subsets.

An Equivalence Relation Partitions a Set

Theorem 2: Let R be an equivalence relation on a set S. Then the equivalence classes of R form a partition of S. Conversely, given a partition $\{A_i \mid i \in I\}$ of the set S, there is an equivalence relation R that has the sets A_i , $i \in I$, as its equivalence classes.

An Equivalence Relation Partitions a Set

Proof: We have already shown the first part of the theorem.

For the second part, assume that $\{A_i \mid i \in I\}$ is a partition of S. Let R be the relation on S consisting of the pairs (x, y) where x and y belong to the same subset A_i in the partition. We must show that R satisfies the properties of an equivalence relation.

- *Reflexivity*: For every $a \in S$, $(a,a) \in R$, because a is in the same subset as itself.
- *Symmetry*: If $(a,b) \in R$, then b and a are in the same subset of the partition, and so is $(b,a) \in R$.
- Transitivity: If $(a,b) \in R$ and $(b,c) \in R$, then a and b are in the same subset of the partition, as are b and c. Since the subsets are disjoint and b belongs to both, the two subsets of the partition must be identical. Therefore, $(a,c) \in R$ since a and c belong to the same subset of the partition. (QED)