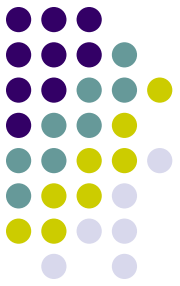
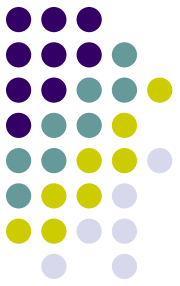


# Type system implementation



- We extend our simple3 language to simple4 with the addition of a type system with three types:
  - int
  - float
  - string
- We also assume that int is a subtype of float and float is a subtype of string, that is, a compiler/interpreter is allowed to insert widening conversions and should flag errors for narrowing conversions.

# Type system implementation



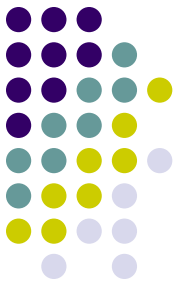
- We want to be able to write programs such as these:

```
int inc(int x) return x+1;
int y = inc(3);
put "the result is", y;
```

```
float pow(float b,int p) {
  if (p == 0)
    return 1.0;
  else
    return b*pow(b,p-1);
}

float v;
get v;
int p;
get p;
float result = pow(v,p);
put v," to the power of ",p," is ",result;
```

# Type system implementation: Syntax



```
prog      :      stmt+;

stmt      :      dataType VAR '(' formalParamList? ')' stmt // declare a function
           |      dataType VAR ('=' exp)? ';' // declare variable in current scope with optional initializer
           |      VAR '=' exp ';' // assign value to variable
           |      'get' (prompt ',')? VAR ';' // prompt user for a value and assign it to variable
           |      'put' exp (',' exp)* ';' // print out value(s) to terminal
           |      VAR '(' actualParamList? ')' ';' // function call statement
           |      'return' exp? ';'
           |      'while' '(' exp ')' stmt
           |      'if' '(' exp ')' stmt ('else' stmt)?
           |      '{' stmt+ '}' // block statement (new local scope)
           ;

dataType  :      'int'
           |      'float'
           |      'string'
           ;

formalParamList
           :      dataType VAR (',' dataType VAR)*
           ;

actualParamList
           :      exp (',' exp)*
           ;
```

# Type system implementation: Syntax

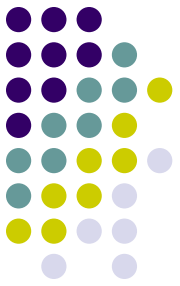


```
prompt      :          string;

exp         :          relexp;
relexp      :          addexp (('==' addexp) | ('<=' addexp))*;
addexp      :          mulexp (('+' mulexp) | ('-' mulexp))*;
mulexp      :          atom (('*' atom) | ('/' atom))*;

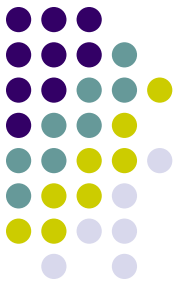
atom        :          '(' exp ')'
                |          VAR '(' actualParamList? ')' // function call within an expression
                |          VAR
                |          '-? INT
                |          '-? FLOAT
                |          string
                ;
```

# Type system implementation: Semantics



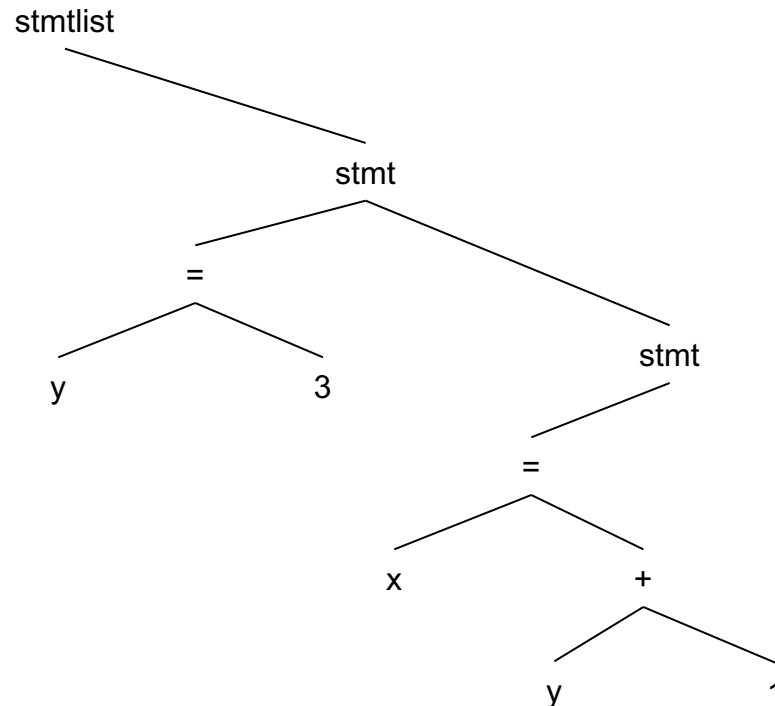
- At the semantic level we *annotate* all ASTs with type information
- We use *type propagation* to check that expressions/statements are properly typed.
  - Type propagation is the systematic tagging of an AST from leafs up with type information.

# Type system implementation: Semantics

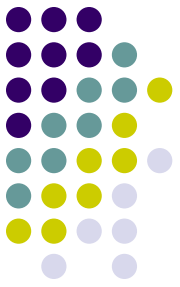


- Consider the simple example:

```
int y;  
int x;  
y = 3;  
x = y + 1;
```

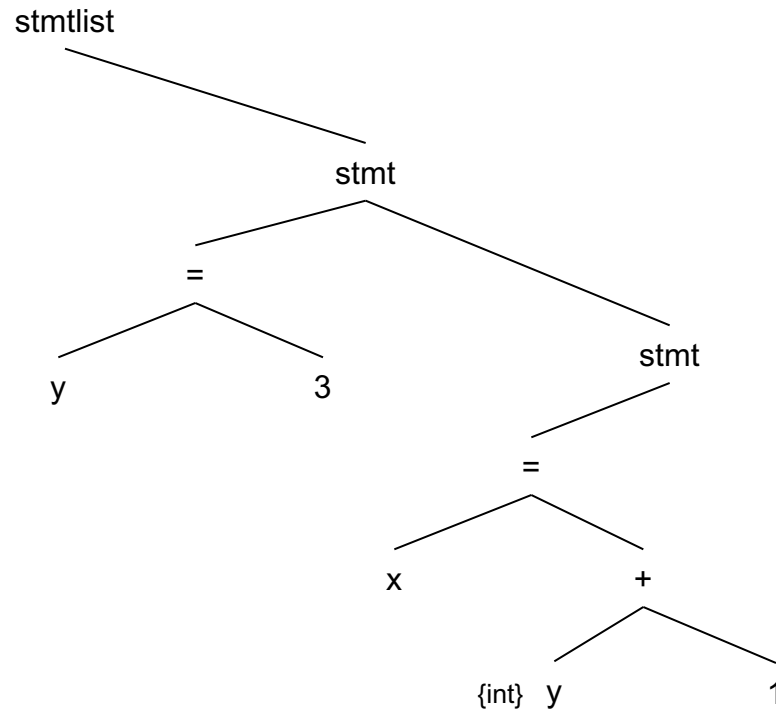


# Type system implementation: Semantics



- Consider the simple example:

```
int y;  
int x;  
y = 3;  
x = y + 1;
```

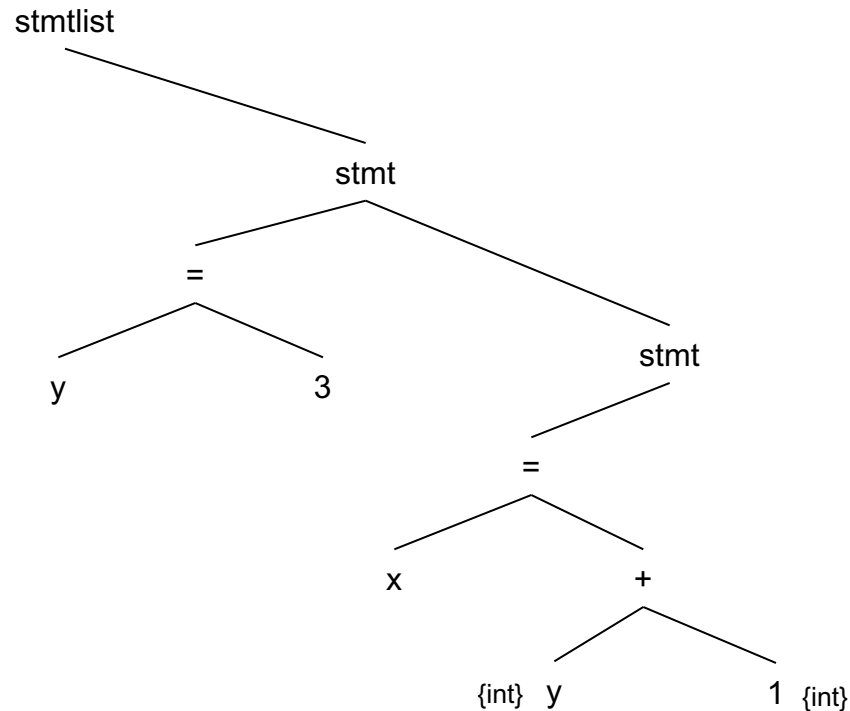


# Type system implementation: Semantics



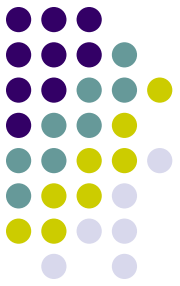
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```
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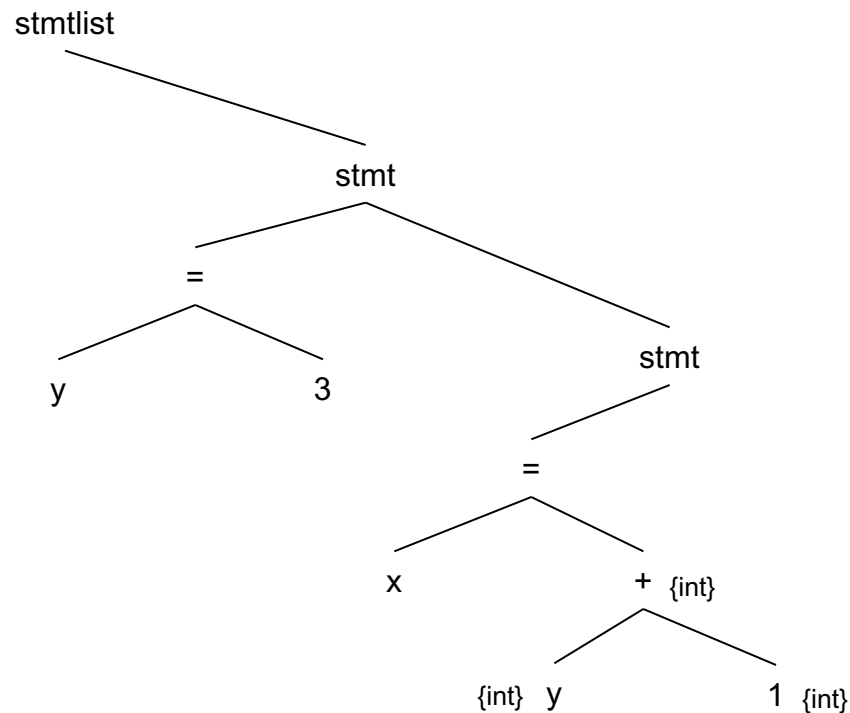


# Type system implementation: Semantics

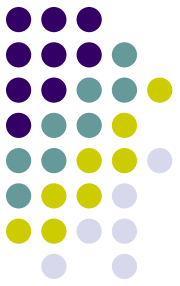


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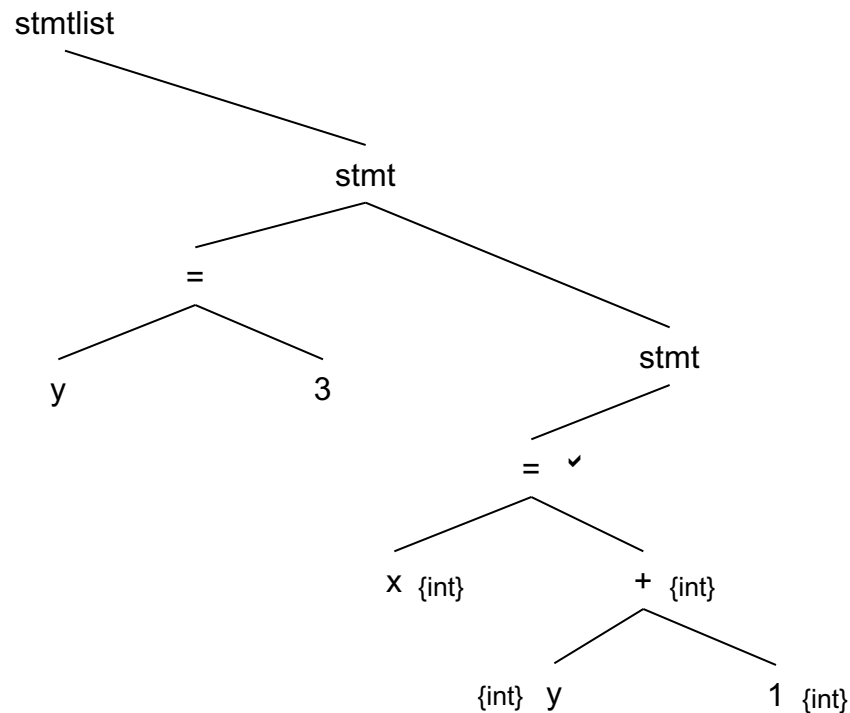


# Type system implementation: Semantics

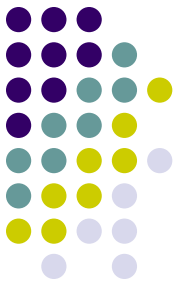


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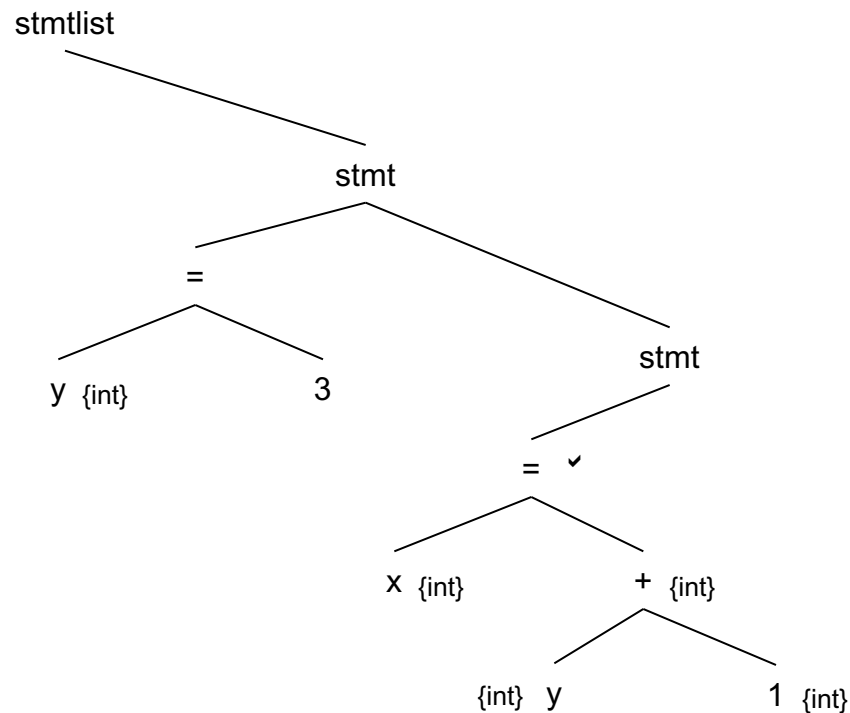


# Type system implementation: Semantics



- Consider the simple example:

```
int y;  
int x;  
y = 3;  
x = y + 1;
```

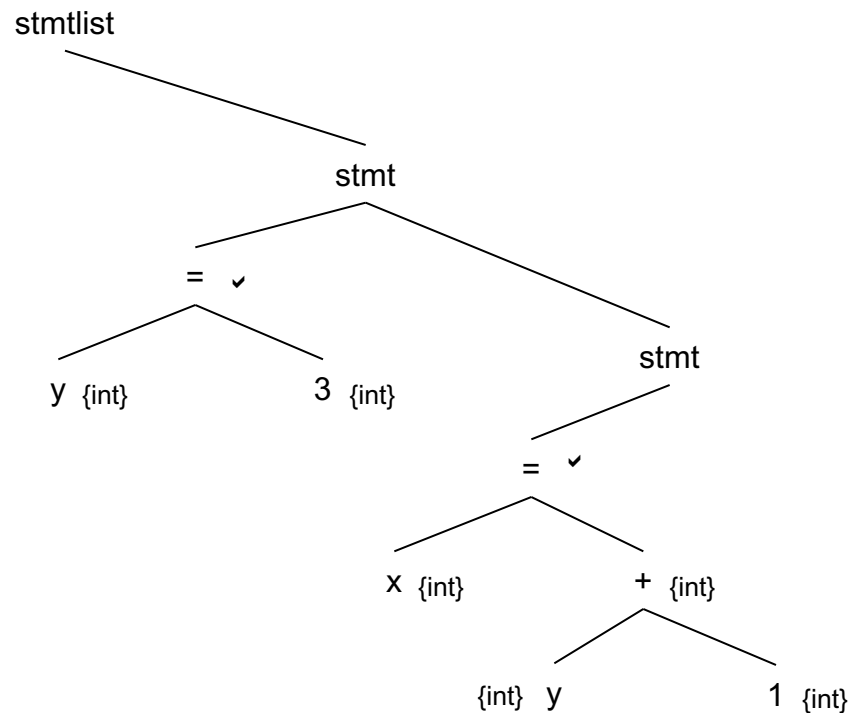


# Type system implementation: Semantics

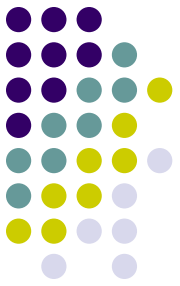


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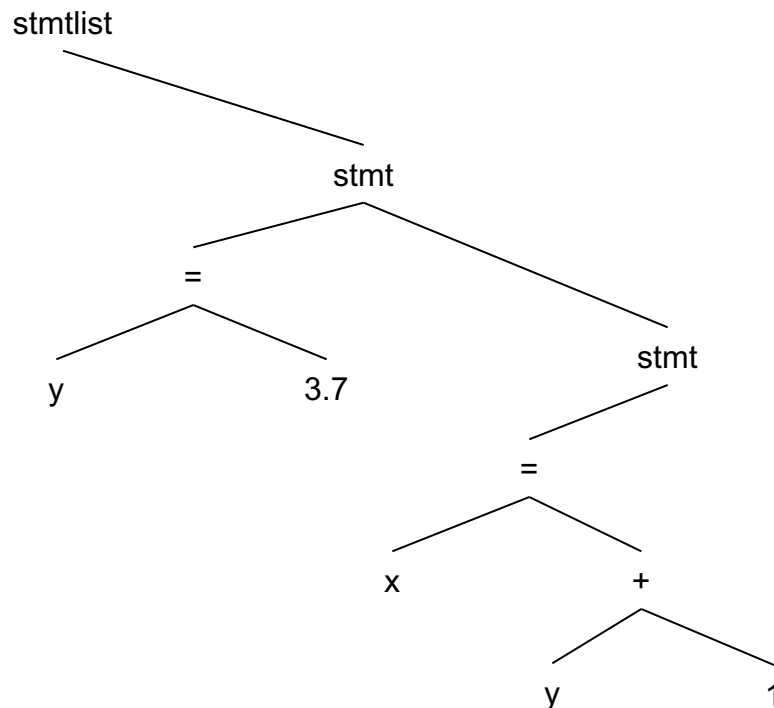


# Type system implementation: Semantics

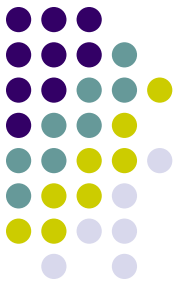


- Consider this example which has a typecheck error:

```
float y;  
int x;  
y = 3.7;  
x = y + 1;
```

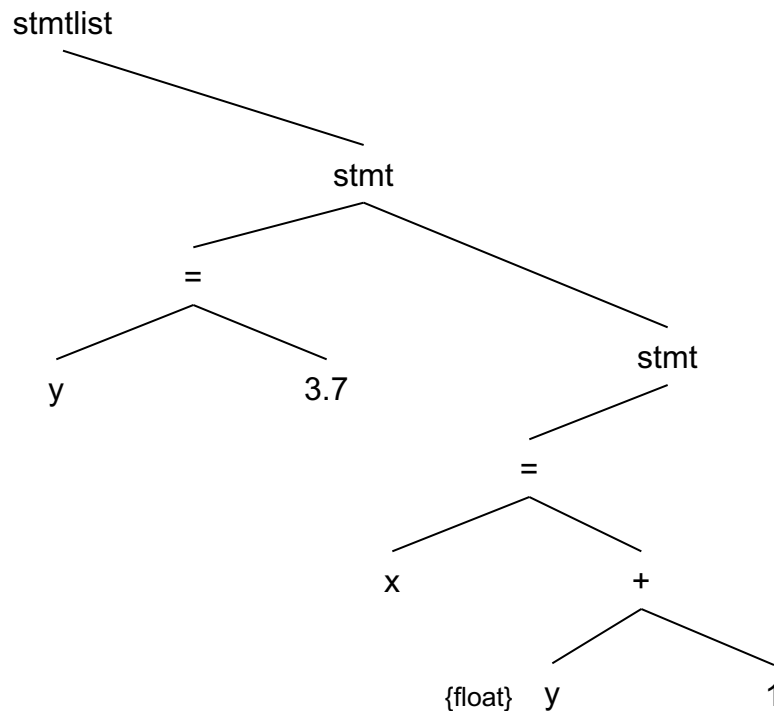


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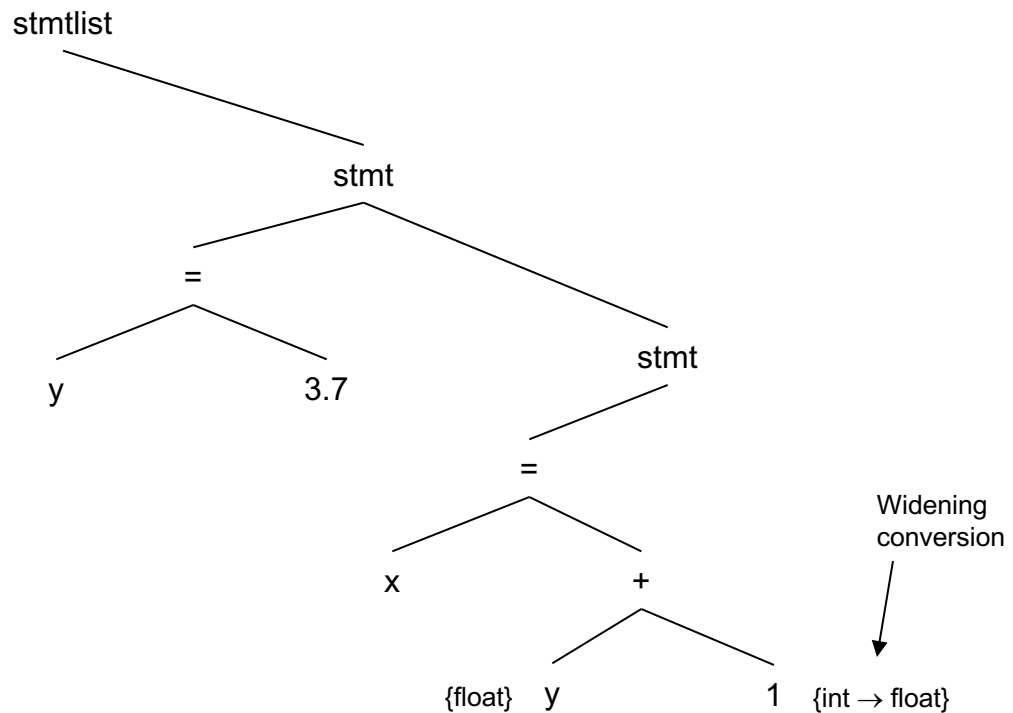


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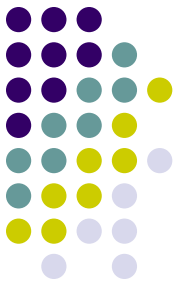


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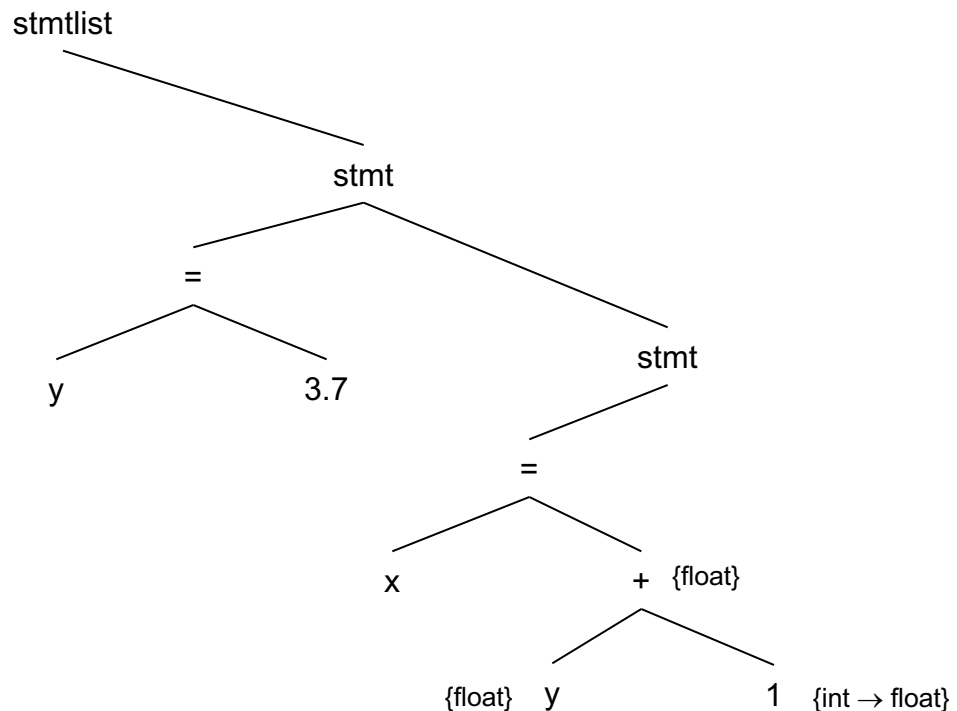


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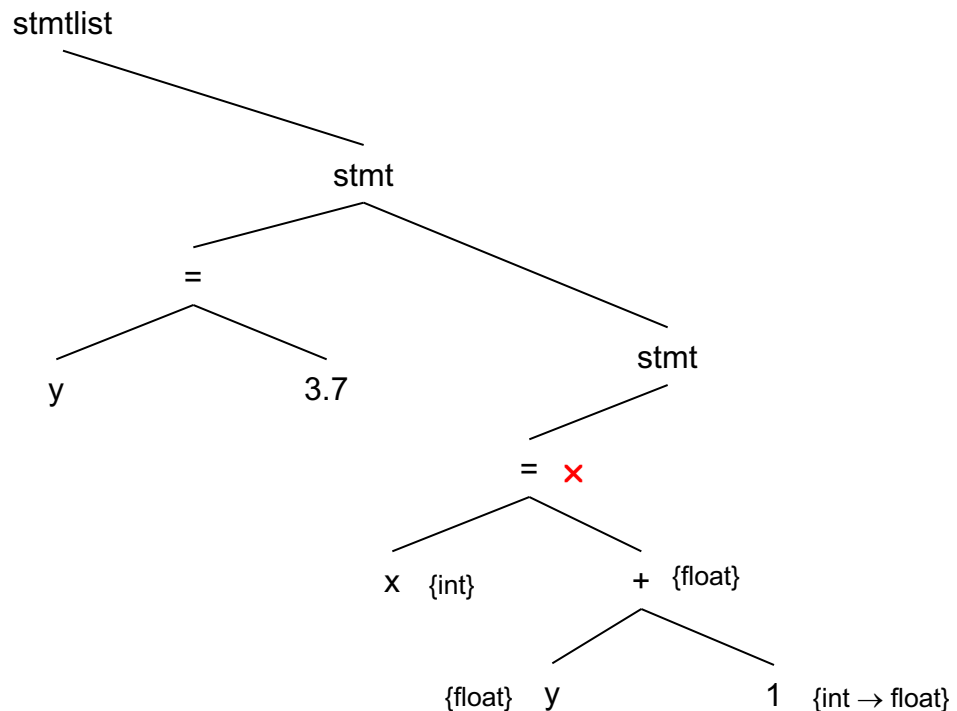


# Type system implementation: Semantics

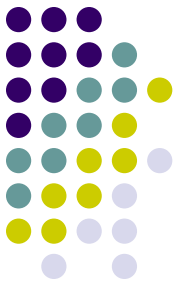


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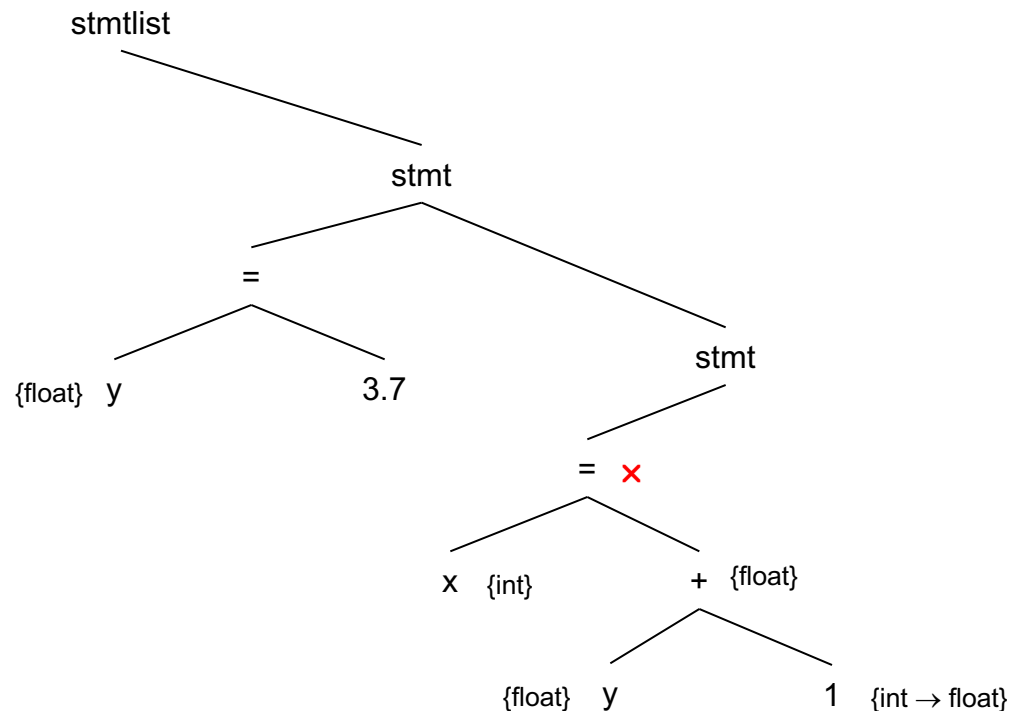


# Type system implementation: Semantics

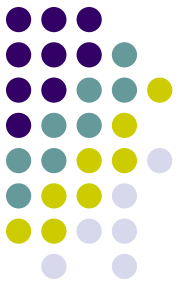


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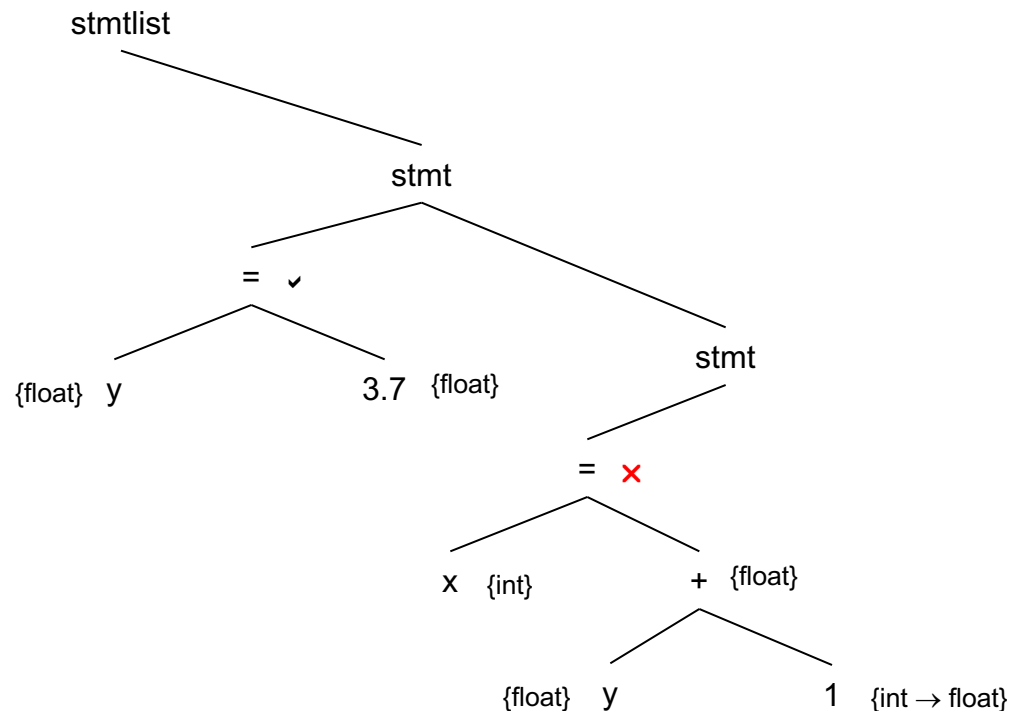


# Type system implementation: Semantics



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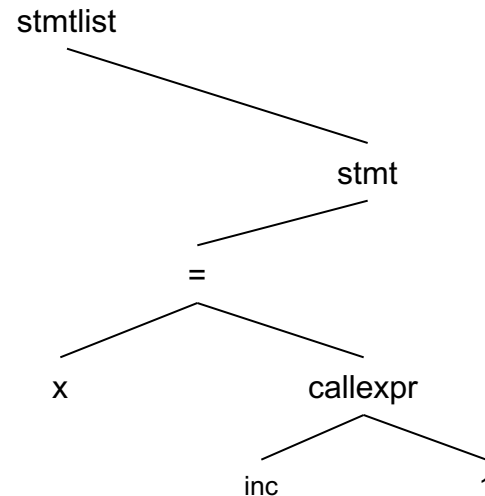


# Type system implementation: Semantics

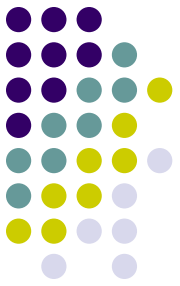


- Here is an example with a function call:

```
int inc(int i) return i+1;  
int x;  
x = inc(1);
```

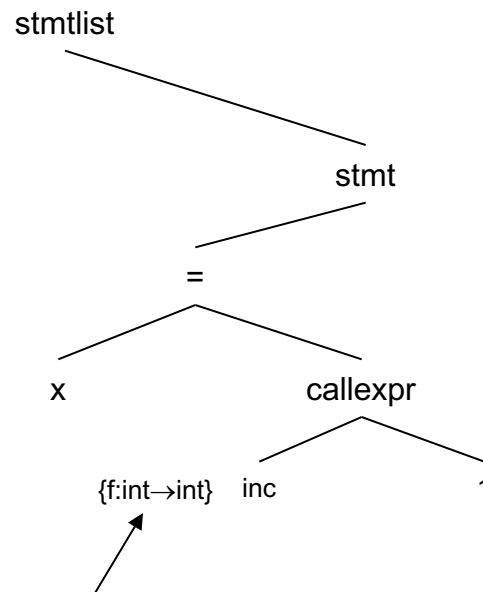


# Type system implementation: Semantics



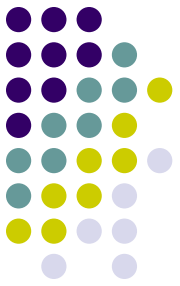
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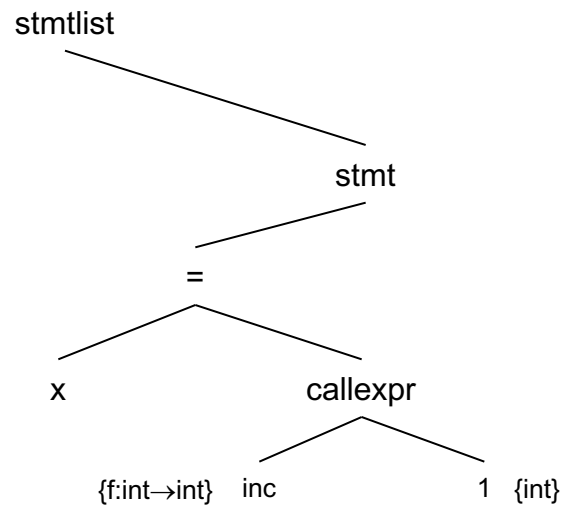
We have to track function symbols, both for their formal parameter types and return types.

# Type system implementation: Semantics

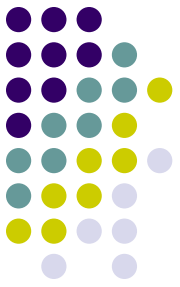


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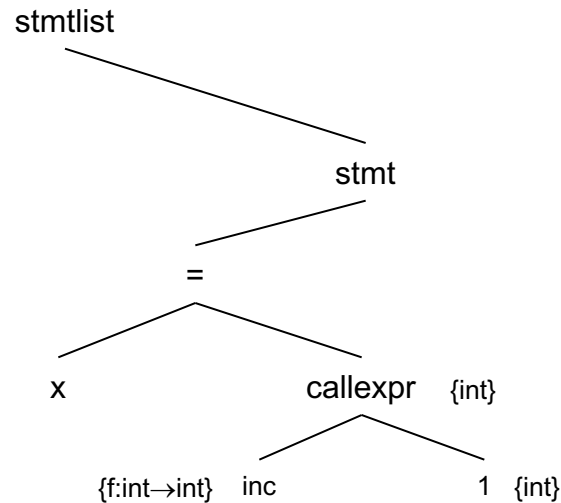


# Type system implementation: Semantics

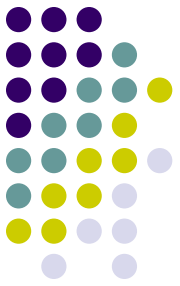


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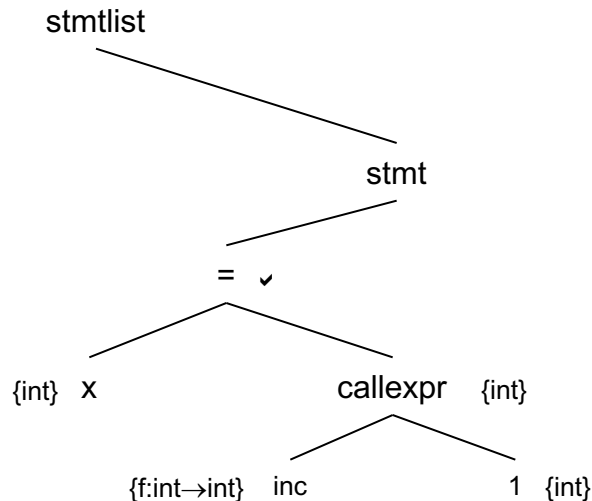


# Type system implementation: Semantics



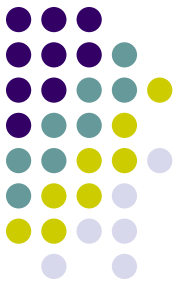
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```
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x = inc(1);
```



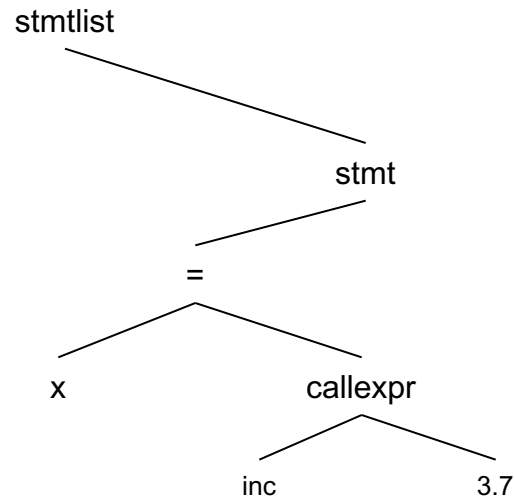


# Type system implementation: Semantics

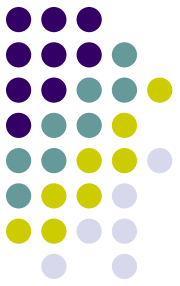


- Here is an example with a function call and a type error:

```
int inc(int i) return i+1;  
int x;  
x = inc(3.7);
```

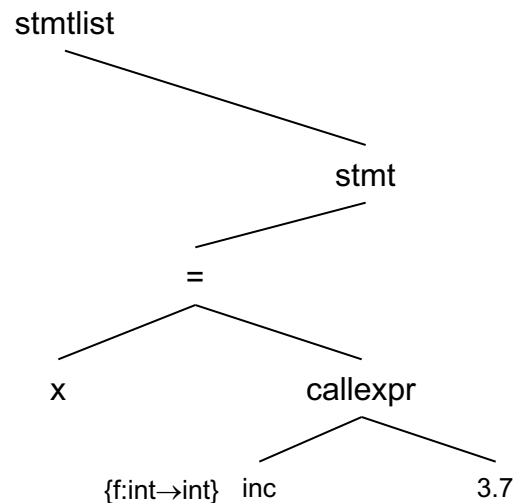


# Type system implementation: Semantics

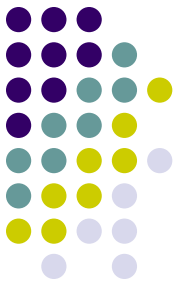


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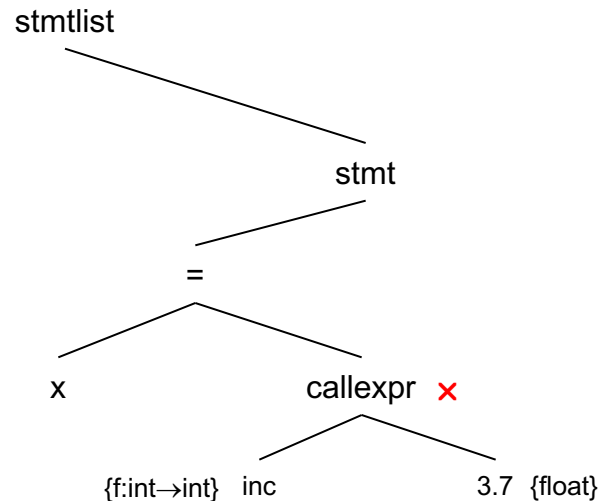


# Type system implementation: Semantics



- Here is an example with a function call and a type error:

```
int inc(int i) return i+1;  
int x;  
x = inc(3.7);
```

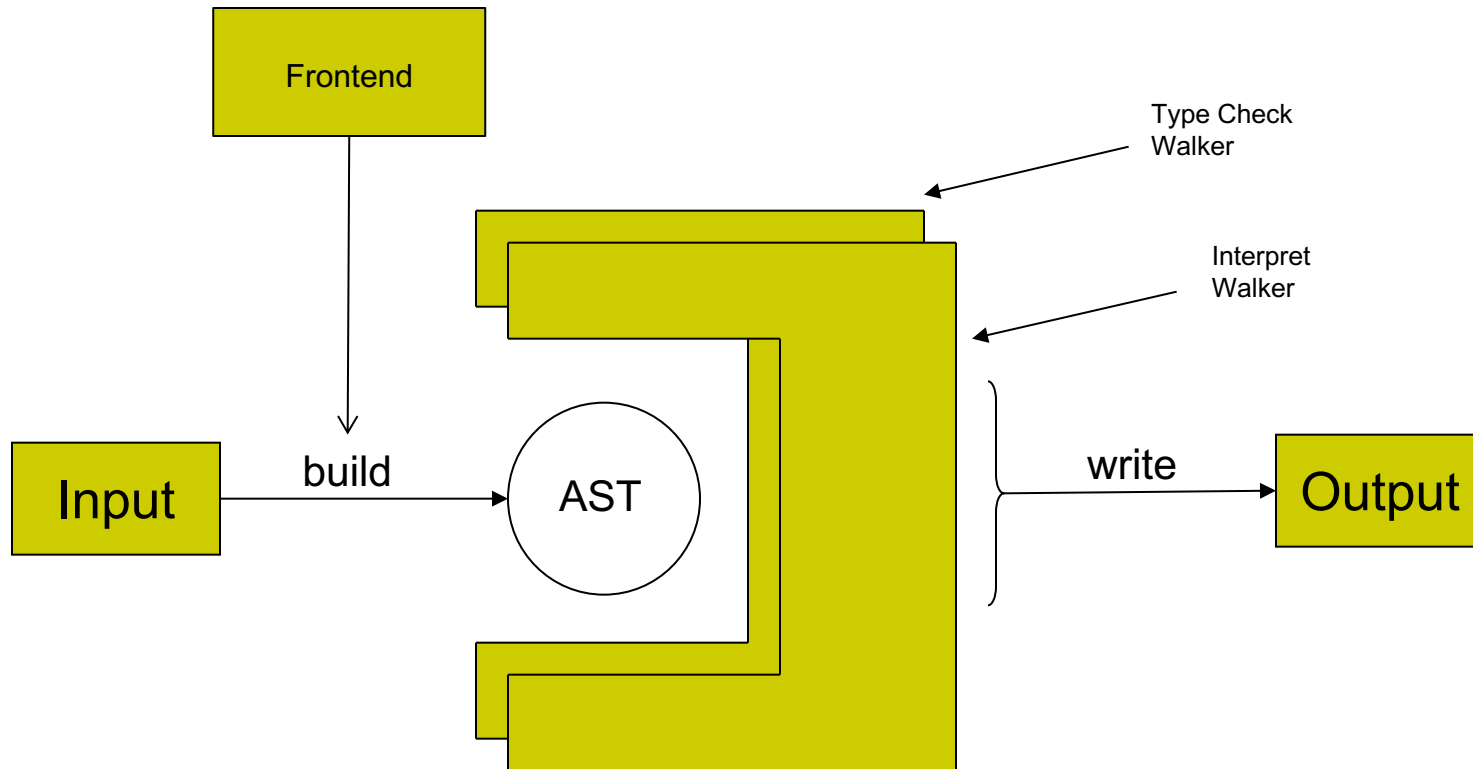
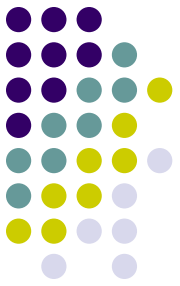


# Type system implementation

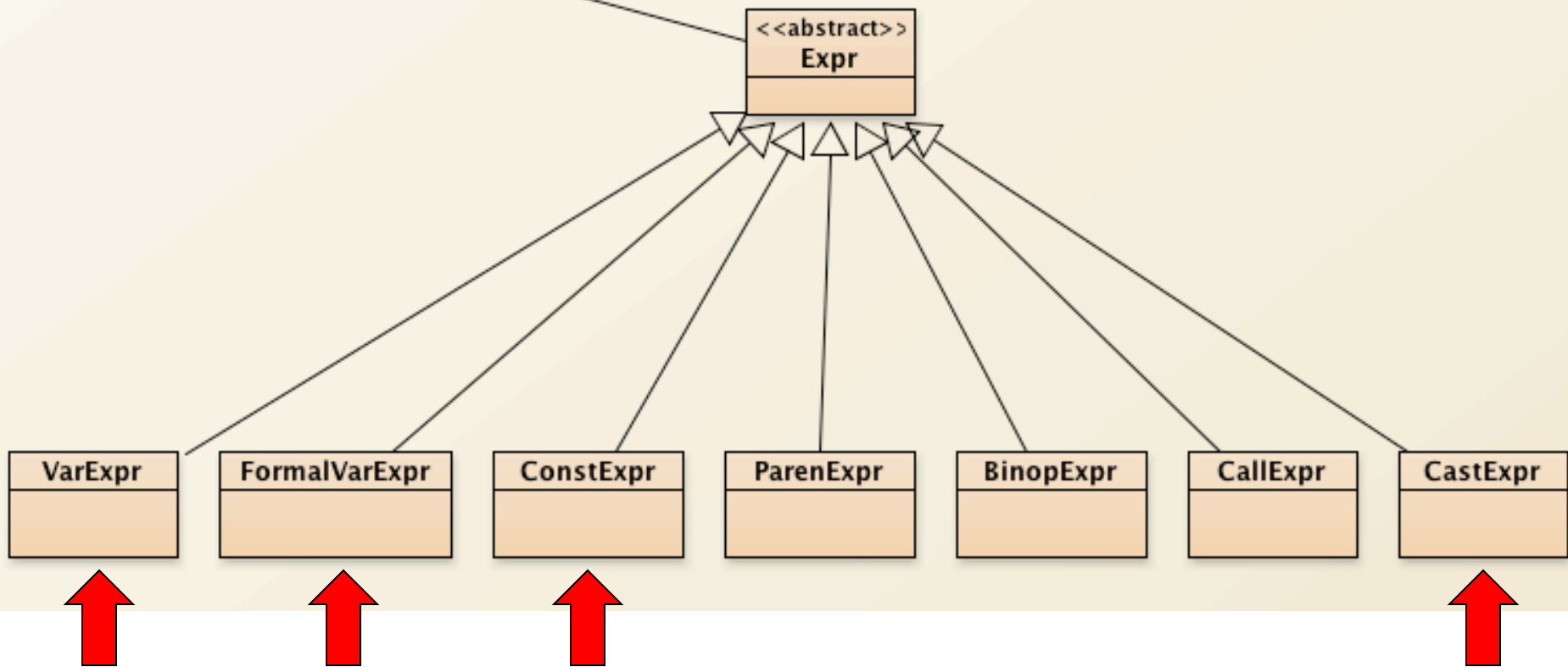
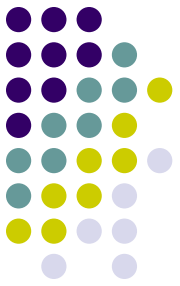


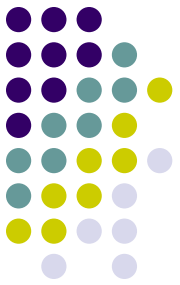
- To implement the type system we introduce the following type tags in the interpreter / symbol table
  - int
  - float
  - string
  - function
- We implement a *static type checker* as a separate walker in the interpreter

# Architecture of our Interpreter



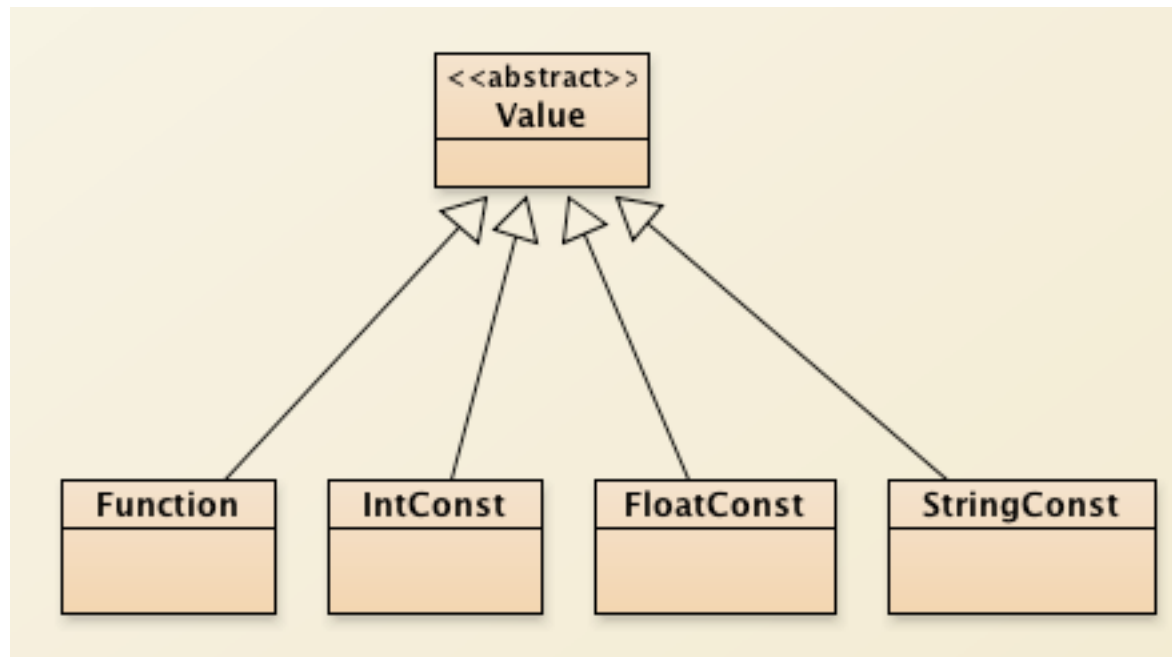
# Extended AST





# Extended AST

- In our previous implementations we only had function values and integer scalars
- We now have additional value types



# The Value Base Class



```
abstract class Value {
    public static final int NOTYPE = -1;
    public static final int INTEGER = 0;
    public static final int FLOAT = 1;
    public static final int STRING = 2;
    public static final int FUNCTION = 3;

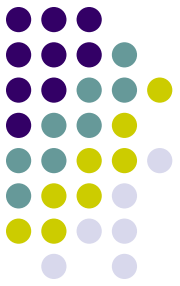
    // Type Promotion Table
    // This table implements the following type hierarchy:
    // int < float < string
    // Note: functions are not allowed to appear
    // in the context of any operations.
    private static int[][] typeArray = {
        // INTEGER FLOAT STRING FUNCTION
        //-----
        { INTEGER, FLOAT, STRING, NOTYPE }, // INTEGER
        { FLOAT, FLOAT, STRING, NOTYPE }, // FLOAT
        { STRING, STRING, STRING, NOTYPE }, // STRING
        { NOTYPE, NOTYPE, NOTYPE, NOTYPE } // FUNCTION
    };

    public static int getResultType(int lt,int rt) {
        if (lt == NOTYPE || rt == NOTYPE)
            return NOTYPE;
        else
            return typeArray[lt][rt];
    }

    // every derived class needs to implement the following behavior
    public abstract int getType();
    public abstract String toString();
}
```



# Value Classes



```
class IntConst extends Value {  
    private Integer value;  
  
    public IntConst(String value) {  
        this.value=new Integer(Integer.parseInt(value));  
    }  
  
    public IntConst(Integer value) {  
        this.value=value;  
    }  
  
    public Integer getValue() {  
        return value;  
    }  
  
    public int getType() {  
        return Value.INTEGER;  
    }  
  
    public String toString() {  
        return value.toString();  
    }  
}
```

```
class StringConst extends Value {  
    private String value;  
  
    public StringConst(String value) {  
        this.value=value;  
    }  
  
    public String getValue() {  
        return value;  
    }  
  
    public int getType() {  
        return Value.STRING;  
    }  
  
    public String toString() {  
        return value;  
    }  
}
```

# Symbol Table

```
public class SymbolTableScope {
    // scope stack is built as a linked list
    private SymbolTableScope parentScope = null;

    // function value, if this is the scope of a function call
    // otherwise null
    Function function = null;

    // symbols are kept in a hashmap indexed by their name
    // their initialization value depends on their kind:
    //   integer,float, string constants, function values
    //
    private HashMap<String, Value> value = new HashMap<String, Value>();

    public SymbolTableScope(SymbolTableScope parentScope) {
        this.parentScope = parentScope;
    }

    public SymbolTableScope getParentScope() {
        return parentScope;
    }

    public void setParentScope(SymbolTableScope parentScope) {
        this.parentScope = parentScope;
    }

    public void enterSymbol(String name, Value value) {
        this.value.put(name, value);
    }

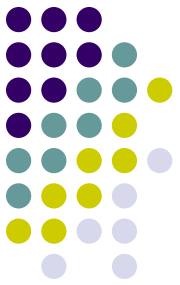
    public Value lookupSymbol(String name) {
        return value.get(name);
    }

    public void setFunctionValue(Function value) {
        function = value;
    }

    public Function getFunctionValue() {
        return function;
    }
}
```



# The Reader



```
stmt returns [Stmt ast]
  // function declarations can have a void return type...
: 'void' VAR '(' ')' s=stmt
  { $ast = new FuncDeclStmt($VAR.text,new Function(Value.NOTYPE,new ArgList(),$s.ast)); }
| dt=dataType VAR '(' ')' s=stmt
  { $ast = new FuncDeclStmt($VAR.text,new Function($dt.type,new ArgList(),$s.ast)); }
| 'void' VAR '(' l=formalParamList ')' s=stmt
  { $ast = new FuncDeclStmt($VAR.text,new Function(Value.NOTYPE,$l.ast,$s.ast)); }
| dt=dataType VAR '(' l=formalParamList ')' s=stmt
  { $ast = new FuncDeclStmt($VAR.text,new Function($dt.type,$l.ast,$s.ast)); }
| dt=dataType VAR '=' exp ';'
  { $ast = new VarDeclStmt($dt.type,$VAR.text,$exp.ast); }
| dt=dataType VAR ';'
  { $ast = new VarDeclStmt($dt.type,$VAR.text,new ConstExpr(new IntConst("0"))); }
| VAR '=' exp ';'
  { $ast = new AssignStmt($VAR.text,$exp.ast); }
| 'get' prompt ',' VAR ';'
  { $ast = new GetStmt($prompt.text,$VAR.text); }
| 'get' VAR ';'
  { $ast = new GetStmt("", $VAR.text); }
| 'put' argList ';'
  { $ast = new PutStmt($argList.ast); }
| VAR '(' l=actualParamList ')' ';'
  { $ast = new CallStmt($VAR.text,$l.ast);}
```

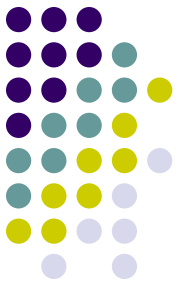
...

```
dataType returns [int type]
: 'int'    {$type=Value.INTEGER;}
| 'float'  {$type=Value.FLOAT;}
| 'string' {$type=Value.STRING;}
;
```

# The Reader



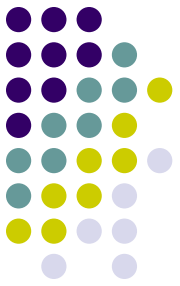
```
atom returns [Expr ast]
: (' exp ')          { $ast = new ParenExpr($exp.ast); }
| VAR (' l=actualParamList ') { $ast = new CallExpr($VAR.text,$l.ast);}
| VAR (' ')          { $ast = new CallExpr($VAR.text);}
| VAR                { $ast = new VarExpr($VAR.text); }
| '-' INT            { $ast = new ConstExpr(new IntConst('-'+$INT.text)); }
| INT                { $ast = new ConstExpr(new IntConst($INT.text)); }
| '-' FLOAT          { $ast = new ConstExpr(new FloatConst('-'+$FLOAT.text)); }
| FLOAT              { $ast = new ConstExpr(new FloatConst($FLOAT.text)); }
| string              { $ast = new ConstExpr(new StringConst($string.text)); }
;
```



# The Type Check Walker

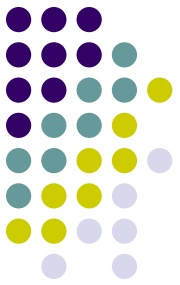
- The type check walker looks like an interpreter...
- ...but it computes types instead of values.

# The Type Check Visitor



```
// the dispatcher for the type check visitor - returns a type tag
public int dispatch(AST ast) {
    if (ast.getClass() == AssignStmt.class) return interp((AssignStmt)ast);
    else if (ast.getClass() == BlockStmt.class) return interp((BlockStmt)ast);
    else if (ast.getClass() == GetStmt.class) return interp((GetStmt)ast);
    else if (ast.getClass() == IfStmt.class) return interp((IfStmt)ast);
    else if (ast.getClass() == PutStmt.class) return interp((PutStmt)ast);
    else if (ast.getClass() == WhileStmt.class) return interp((WhileStmt)ast);
    else if (ast.getClass() == StmtList.class) return interp((StmtList)ast);
    else if (ast.getClass() == BinopExpr.class) return interp((BinopExpr)ast);
    else if (ast.getClass() == ConstExpr.class) return interp((ConstExpr)ast);
    else if (ast.getClass() == ParenExpr.class) return interp((ParenExpr)ast);
    else if (ast.getClass() == VarExpr.class) return interp((VarExpr)ast);
    else if (ast.getClass() == FuncDeclStmt.class) return interp((FuncDeclStmt)ast);
    else if (ast.getClass() == VarDeclStmt.class) return interp((VarDeclStmt)ast);
    else if (ast.getClass() == CallStmt.class) return interp((CallStmt)ast);
    else if (ast.getClass() == CallExpr.class) return interp((CallExpr)ast);
    else if (ast.getClass() == ReturnStmt.class) return interp((ReturnStmt)ast);
    else {
        System.err.println("Error (InterpVisitor): unknown class type");
        System.exit(1);
        return Value.NOTYPE;
    }
}
```

# The Type Check Visitor



```
// assignment statements
private int interp(AssignStmt ast) {

    // typecheck the expression
    int exprType = this.dispatch(ast.getAST(0));
    // get the type of the variable
    int varType = Interpret.symbolTable.lookupSymbol(ast.lhsVar()).getType();

    // types compatible?
    int resultType = Value.getResultType(varType,exprType);

    // check for type errors
    if (resultType == Value.NOTYPE ||
        resultType != varType) { // second condition means: assigning supertype to subtype
        System.err.println("Error (assignmentstmt): expression type "+resultType+" cannot be assigned to variable of type "+varType);
        System.exit(1);
        return Value.NOTYPE;
    }

    // check if we have to insert a type promotion
    if (resultType != exprType) {
        AST newAst = new CastExpr(exprType,resultType,(Expr)ast.getAST(0));
        ast.putAST(0,newAst);
    }

    // statements do not have types
    return Value.NOTYPE;
}
```

# The

```
// binop expressions
private int interp(BinopExpr ast) {

    // typecheck left child
    int leftType = this.dispatch(ast.getAST(0));

    // typecheck right child
    int rightType = this.dispatch(ast.getAST(1));

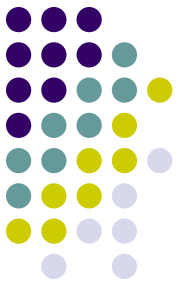
    // see if the expression is well typed
    int resultType = Value.getResultType(leftType,rightType);

    // check for type errors
    // NOTE: add on type string is string concatenation
    if (resultType == Value.NOTYPE) {
        System.err.println("Error (binopexpr): binop expression with types "+leftType+" and "+rightType+" is ill-typed");
        System.exit(1);
        return Value.NOTYPE;
    }

    // check if we have to insert a type promotion
    if (resultType != leftType) {
        AST newAst = new CastExpr(leftType,resultType,(Expr)ast.getAST(0));
        ast.putAST(0,newAst);
    }

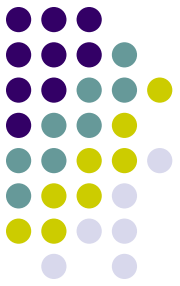
    // check if we have to insert a type promotion
    if (resultType != rightType) {
        AST newAst = new CastExpr(leftType,resultType,(Expr)ast.getAST(1));
        ast.putAST(1,newAst);
    }

    // the result type is correct except for the relational operators which
    // always construct an integer return value.
    if (ast.getOp() == BinopExpr.EQ || ast.getOp() == BinopExpr.LESSEQ) {
        return Value.INTEGER;
    }
    else {
        return resultType;
    }
}
```





# The Type Check Visitor



```
// while statements
private Integer interp(WhileStmt ast) {
    // typecheck the expression -- has to be an integer
    int exprType = this.dispatch(ast.getAST(0));

    if (exprType != Value.INTEGER) {
        System.err.println("Error: expression of a while-stmt has to be of type integer.");
        System.exit(1);
        return Value.NOTYPE;
    }

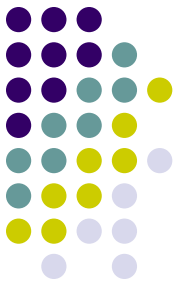
    // type check the body of the loop
    this.dispatch(ast.getAST(1));

    // statements do not have types
    return Value.NOTYPE;
}
```

Note: we do not execute the loop, we simply compute all the types.

# Code

- `SIMPLE4INTERPRETER.zip`



# Assignment

- Assignment #8 – see website

