Grammars and Semantics

- Programming languages are used to specify computations – that is, computations are the meaning/semantics of programs.

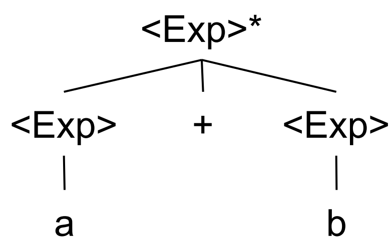
Read Chap 3

Grammars and Semantics

Consider the simple language of expressions:

```
G: <Exp>* ::= <Exp> + <Exp>
           |  <Exp> * <Exp>
           |  a
           |  b
           |  c
```

When we write the sentence `a + b` we can build the parse tree:



We can say that this parse tree *represents* the computation `a + b`.

If we let a and b be variables, then the parse tree gives us a procedure to compute a + b by starting at the leaves of the tree: (1) lookup the values of the variables (2) pass the values up along the parse tree branches (3) use the values to compute the value of the + operator.

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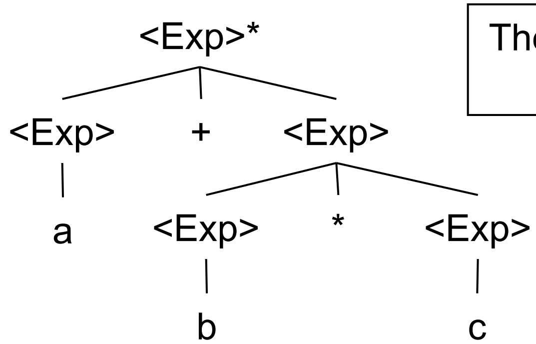
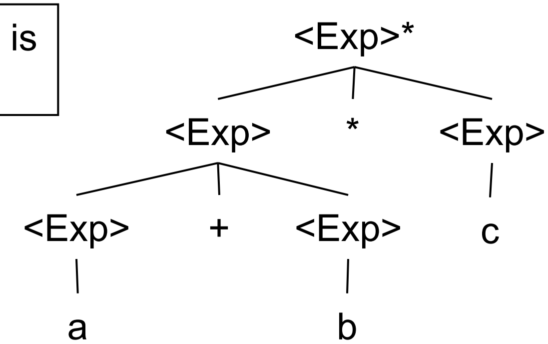
Now consider the sentence `a + b * c`, for this sentence we can construct two parse trees:



The grammar G is ambiguous



Even though both parse trees derive the same terminal string, the computations they represent are very different:

(1) left tree – first compute the product, then the addition
(2) right tree – first compute the addition, then the product

Since we had written the original sentence without parentheses the left parse tree represents the intended computation according to algebraic conventions.

However, from a machine point of view, there is no way of knowing which parse tree to pick…

Grammars and Semantics

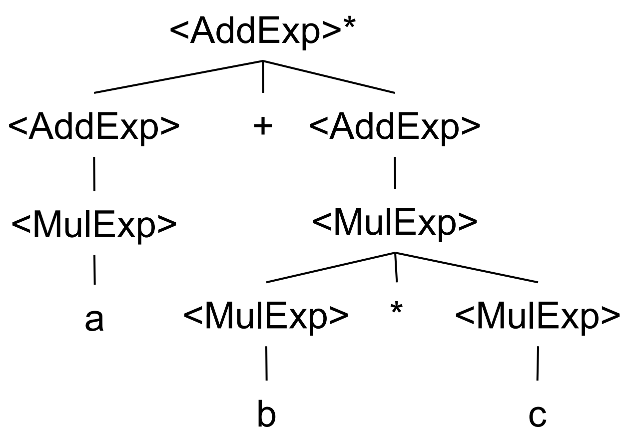
…we need additional information: operator precedence

Operator precedence means that some operators bind tighter than others, e.g. * binds tighter than +.

We can build operator precedence right into our grammar:

```
G': <AddExp>*::= <AddExp> + <AddExp>
             |   <MulExp>
    <MulExp> ::= <MulExpr> * <MulExp>
             | a | b | c
```

Let's try our problematic sentence `a + b * c`, only one parse tree is possible:



This is the only parse tree we can build, therefore, the grammar G' is not ambiguous.

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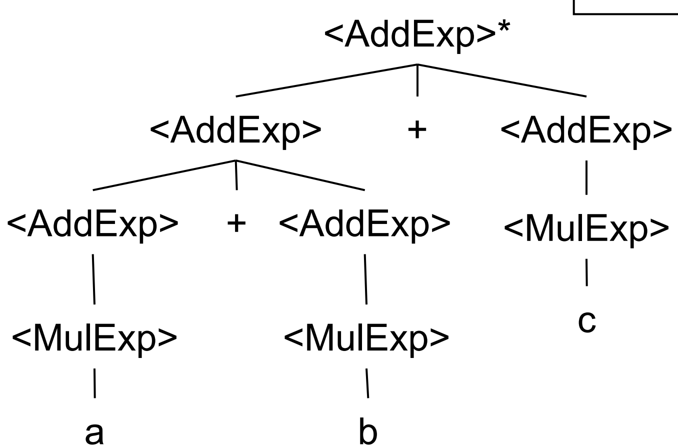
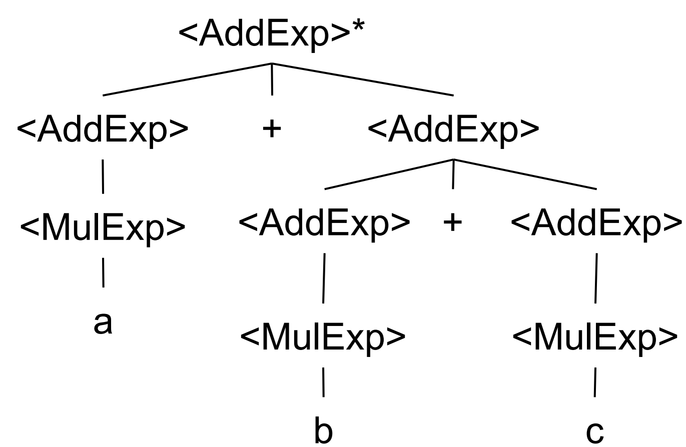
However, our new grammar still has a problem, consider the sentence a+b+c; here we have two possible parse trees:

```
G': <AddExp> ::= <AddExp> + <AddExp>
               |   <MulExp>
     <MulExp> ::= <MulExpr> * <MulExp>
               | a | b | c
```

The grammar G' is ambiguous





Grammars and Semantics

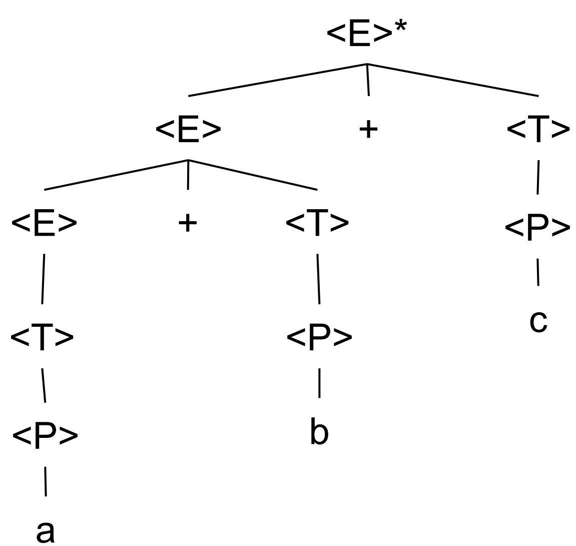
- Again, our grammar is ambiguous because the computation specified by the sentence a+b+c can be represented by two different parse trees.
- We need more information!
- There is one more algebraic property we have not yet explored – associativity
- Most algebraic operators, including the + operator, are left-associative.
- We can rewrite our grammar to take advantage of this additional information:

```
G”: <E>* ::= <E> + <T> | <T>
    <T>  ::= <T> * <P> | <P>
    <P>  ::= a | b | c
```

Grammars and Semantics

Let’s try our sentence a+b+c again with grammar G’’:

```
G”: <E>*  ::= <E> + <T> | <T>
    <T>   ::= <T> * <P> | <P>
    <P>   ::= a | b | c
```



There is no other way to derive this string from the grammar and thus the grammar is not ambiguous.

Take Away

- Grammars can be ambiguous in the sense that a derived string can have multiple distinct parse trees.
- By taking additional information such as associativity and precedence about the operators of a language into account we can construct grammars that are not ambiguous.

Grammars and Semantics

Given the following grammar,

```
G”: <E>* ::= <E> + <T> | <T>
     <T>  ::= <T> * <P> | <P>
     <P>  ::= a | b | c
```

Add productions to the grammar that define the right-associative operator = at a lower precedence than any of the other operators.

This new operator should allow you to write expressions such as

a = b
a = b = c
a = b = b + c

Grammars and Semantics

a) Show that the following grammar is ambiguous.

```
G: <S> ::= <S> <S>
         |    ( <S> )
         |    ()
```

b) Rewrite the above grammar so that it is no longer ambiguous.

Assignment

Assignment #2 – see website